



**Hewlett Packard
Enterprise**

INTRODUCING CHAPEL: A PROGRAMMING LANGUAGE FOR PRODUCTIVE PARALLEL COMPUTING FROM LAPTOPS TO SUPERCOMPUTERS

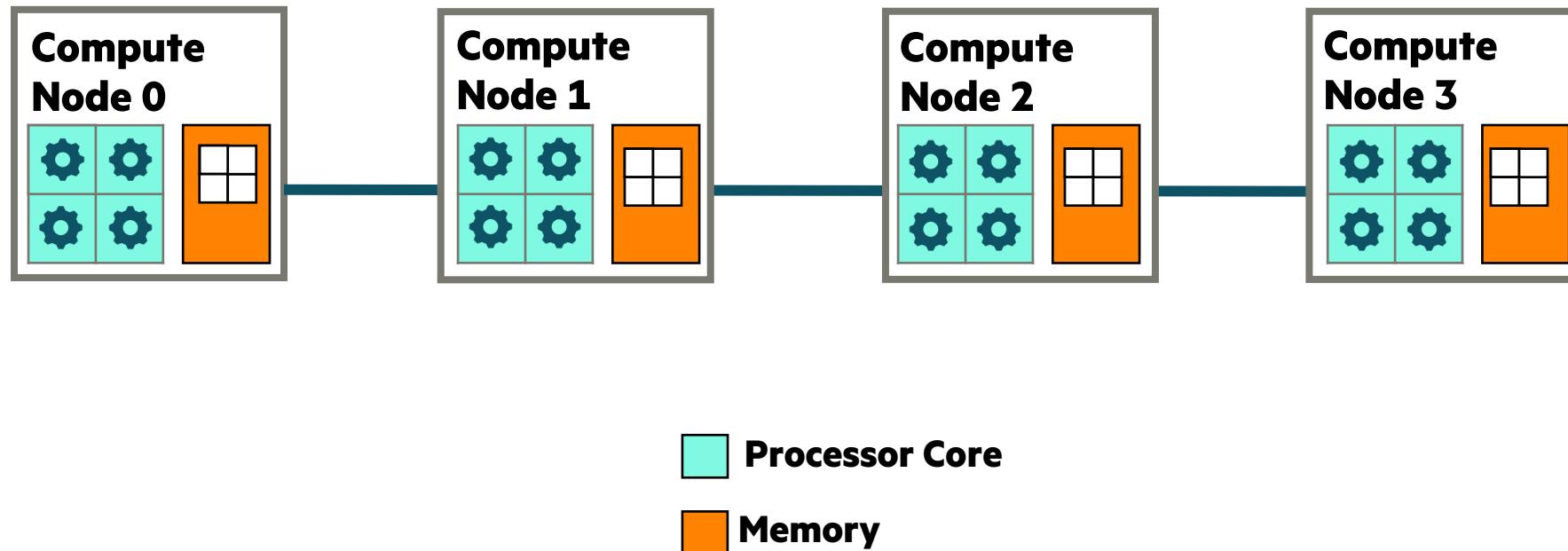
Brad Chamberlain, Distinguished Technologist

LinuxCon, May 11, 2023

PARALLEL COMPUTING IN A NUTSHELL

Parallel Computing: Using the processors and memories of multiple compute resources

- in order to run a program...
 - faster than we could otherwise
 - and/or using larger problem sizes



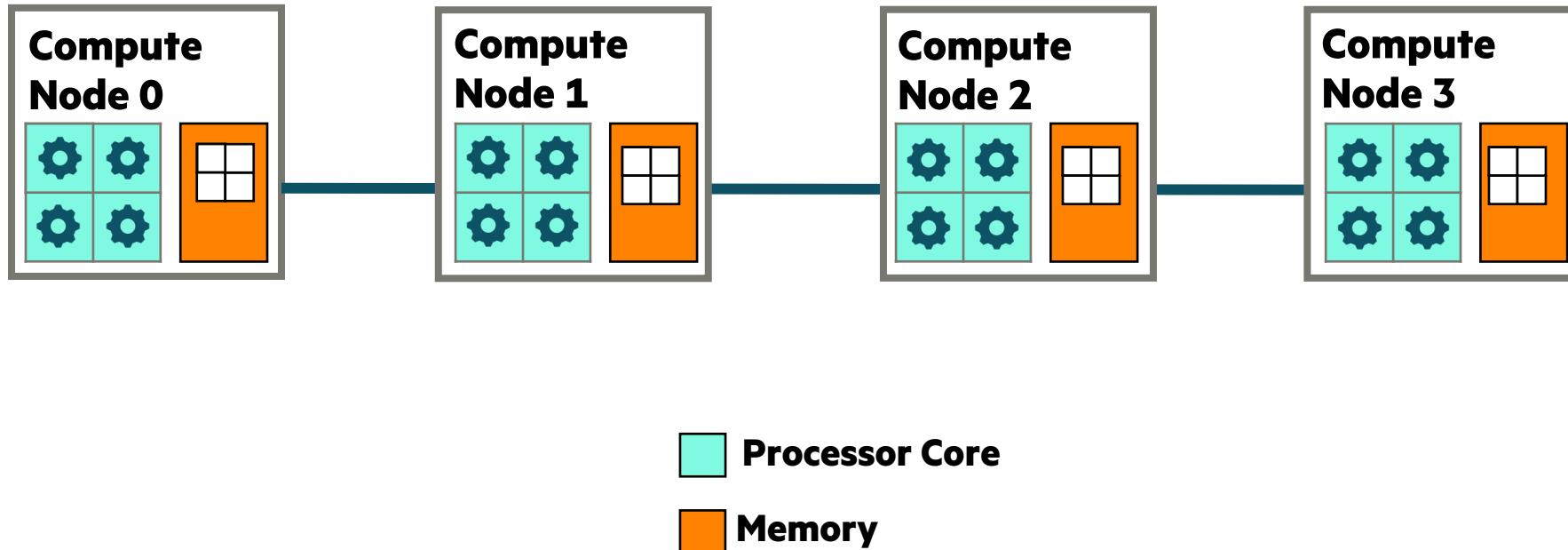
PARALLEL COMPUTING HAS BECOME UBIQUITOUS

Traditional parallel computing:

- supercomputers
- commodity clusters

Today:

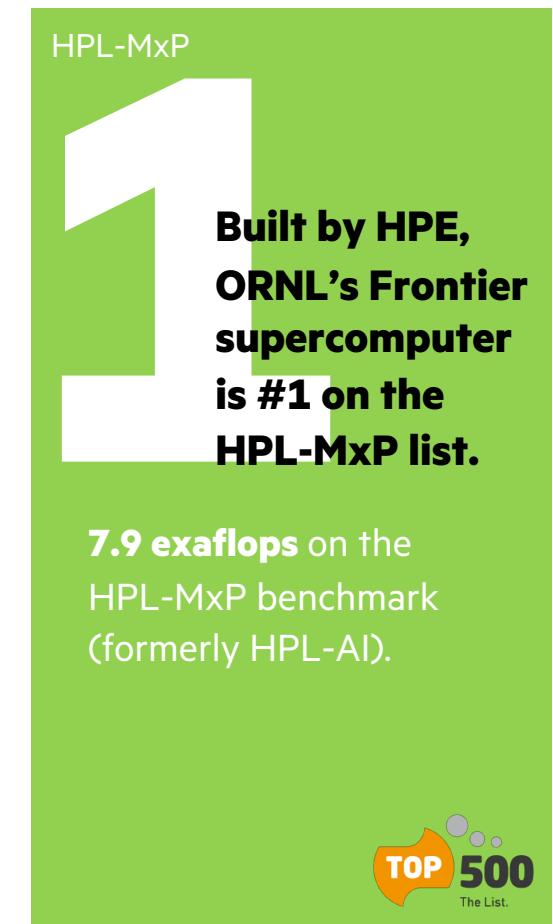
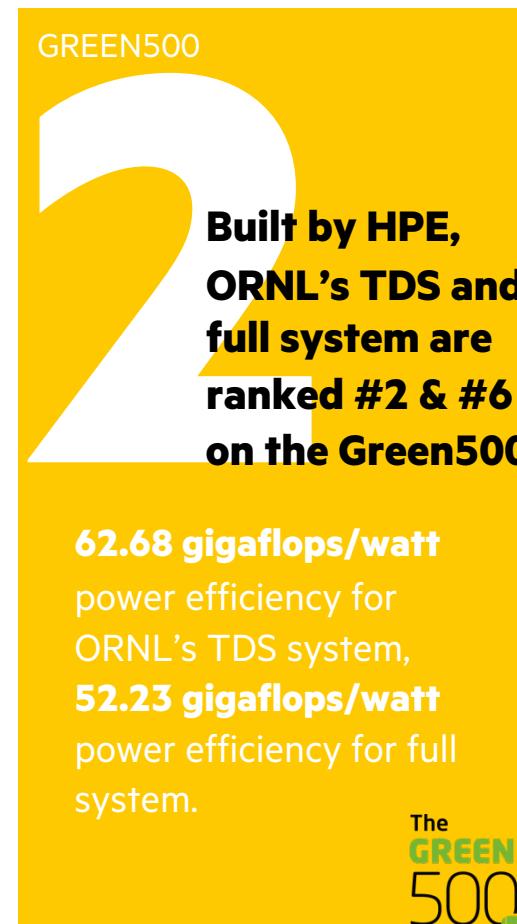
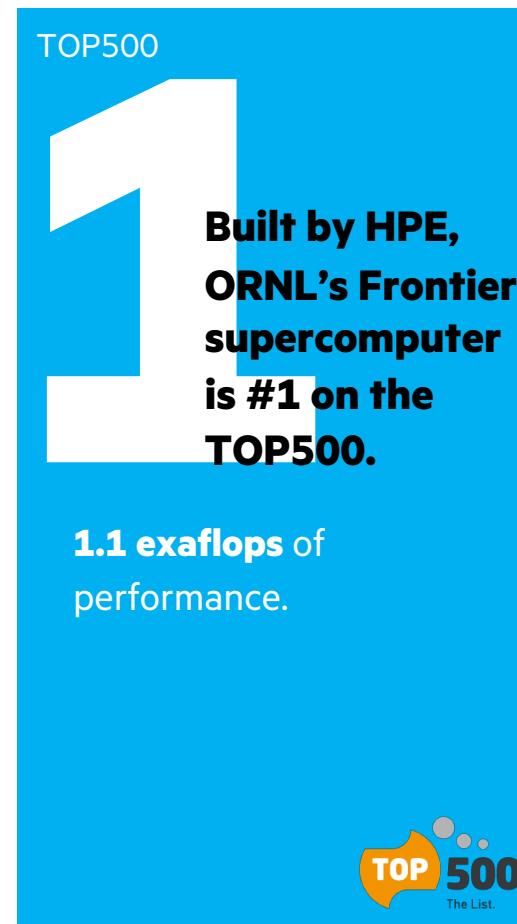
- multicore processors
- GPUs
- cloud computing



OAK RIDGE NATIONAL LABORATORY'S FRONTIER SUPERCOMPUTER



- 74 HPE Cray EX cabinets
- 9,408 AMD CPUs, 37,632 AMD GPUs
- 700 petabytes of storage capacity, peak write speeds of 5 terabytes per second using Cray ClusterStor storage system
- HPE Slingshot networking cables providing 100 GB/s network bandwidth.



Source: May 30, 2022 [Top500](#) release, [HPL-MxP](#) mixed-precision benchmark (formerly HPL-AI).

HPC BENCHMARKS USING CONVENTIONAL PROGRAMMING APPROACHES

STREAM TRIAD: C + MPI + OPENMP

```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Parms *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size(comm, &commSize);
    MPI_Comm_rank(comm, &myRank);

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );

    return errCount;
}

int HPCC_Stream(HPCC_Parms *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3, sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );

    if ( !a || !b || !c ) {
        if ( c ) HPCC_free(c);
        if ( b ) HPCC_free(b);
        if ( a ) HPCC_free(a);
        if ( doIO ) {
            fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
            fclose( outFile );
        }
        return 1;
    }

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++) {
        b[j] = 2.0;
        c[j] = 1.0;
    }
    scalar = 3.0;

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    for (j=0; j<VectorSize; j++) {
        a[j] = b[j]+scalar*c[j];
    }

    HPCC_free(c);
    HPCC_free(b);
    HPCC_free(a);

    return 0;
}
```

HPCC RA: MPI KERNEL

```
/* Perform updates to main table. The scalar equivalent is:
 * for (i=0;i<NUPDATE;i++) {
 *     Ran = (Ran << 1) ^ ((s64)4m) Ran < 0 ? POLY : 0;
 *     Table[Ran & (TABSIZ-1)] = Ran;
 * }
 */

MPI_Irecv(4LocalRecvBuffer, localBufferSize, tparams.dtype64,
          MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, &inreq);
while (i < Sendcnt) {
    /* receive messages */
    do {
        MPI_Test(&inreq, &have_done, &status);
        if (have_done) {
            if (status.MPI_TAG == UPDATE_TAG) {
                MPI_Get_count(&status, tparams.dtype64, &recvUpdates);
                bufferBase = 0;
                for (j=0; j < recvUpdates; j++) {
                    imsg = LocalRecvBuffer(bufferBase+j);
                    LocalOffset = (imsg & (tparams.TableSize - 1)) -
                                 tparams.GlobalStartMyProc;
                    HPCC_Table[LocalOffset] ^= imsg;
                }
            } else if (status.MPI_TAG == FINISHED_TAG) {
                NumberReceiving--;
            } else
                MPI_Abort( MPI_COMM_WORLD, -1 );
            MPI_Irecv(4LocalRecvBuffer, localBufferSize, tparams.dtype64,
                      MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, &inreq);
        }
    } while (have_done && NumberReceiving > 0);
    if (pendingUpdates < maxPendingUpdates) {
        Ran = (Ran << 1) ^ ((s64)4m) Ran < 0 ? POLY : ZERO64B;
        GlobalOffset = Ran & (tparams.TableSize-1);
        if ( GlobalOffset < tparams.Top )
            WhichPe = (GlobalOffset / (tparams.MiniLocalTableSize + 1));
        else
            WhichPe = (GlobalOffset - tparams.Remainder) /
                      tparams.MiniLocalTableSize;
        if (WhichPe == tparams.MyProc) {
            LocalOffset = (Ran & (tparams.TableSize - 1)) -
                         tparams.GlobalStartMyProc;
            HPCC_Table[LocalOffset] ^= Ran;
        }
    }
}

MPI_Test(&outreq, &have_done, MPI_STATUS_IGNORE);
if (have_done) {
    outreq = MPI_REQUEST_NULL;
    pe = HPCC_GetUpdates(Buckets, LocalSendBuffer, localBufferSize,
                          &peUpdates);
    MPI_Isend(4LocalSendBuffer, peUpdates, tparams.dtype64, (int)pe,
              UPDATE_TAG, MPI_COMM_WORLD, &outreq);
    pendingUpdates -= peUpdates;
}
}

/* send remaining updates in buckets */
while (pendingUpdates > 0) {
    /* receive messages */
    do {
        MPI_Test(&inreq, &have_done, &status);
        if (status.MPI_TAG == UPDATE_TAG) {
            MPI_Get_count(&status, tparams.dtype64, &recvUpdates);
            bufferBase = 0;
            for (j=0; j < recvUpdates; j++) {
                imsg = LocalRecvBuffer(bufferBase+j);
                LocalOffset = (imsg & (tparams.TableSize - 1)) -
                             tparams.GlobalStartMyProc;
                HPCC_Table[LocalOffset] ^= imsg;
            }
        } else if (status.MPI_TAG == FINISHED_TAG) {
            /* we got a done message. Thanks for playing.. */
            NumberReceiving--;
        } else
            MPI_Abort( MPI_COMM_WORLD, -1 );
        MPI_Irecv(4LocalRecvBuffer, localBufferSize, tparams.dtype64,
                  MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, &inreq);
    } while (have_done && NumberReceiving > 0);
}

MPI_Waitall( tparams.NumProcs, tparams.finish_red, tparams.finish_statuses);
```

SCALABLE PARALLEL PROGRAMMING THAT'S AS NICE AS PYTHON?

Imagine having a programming language for parallel computing that was as...

...**programmable** as Python

...yet also as...

...**fast** as Fortran/C/C++

...**scalable** as MPI/SHMEM

...**GPU-ready** as CUDA/OpenMP/OpenCL/OpenACC/...

...**portable** as C

...**fun** as [your favorite programming language]

This is our motivation for Chapel



WHAT IS CHAPEL?

Chapel: A modern parallel programming language

- portable & scalable
- open-source & collaborative



Goals:

- Support general parallel programming
- Make parallel programming at scale far more productive



FIVE KEY CHARACTERISTICS OF CHAPEL

1. **compiled:** to generate the best performance possible
2. **statically typed:** to avoid simple errors after hours of execution
3. **interoperable:** with C, Fortran, Python, ...
4. **portable:** runs on laptops, clusters, the cloud, supercomputers
5. **open-source:** to reduce barriers to adoption and leverage community contributions



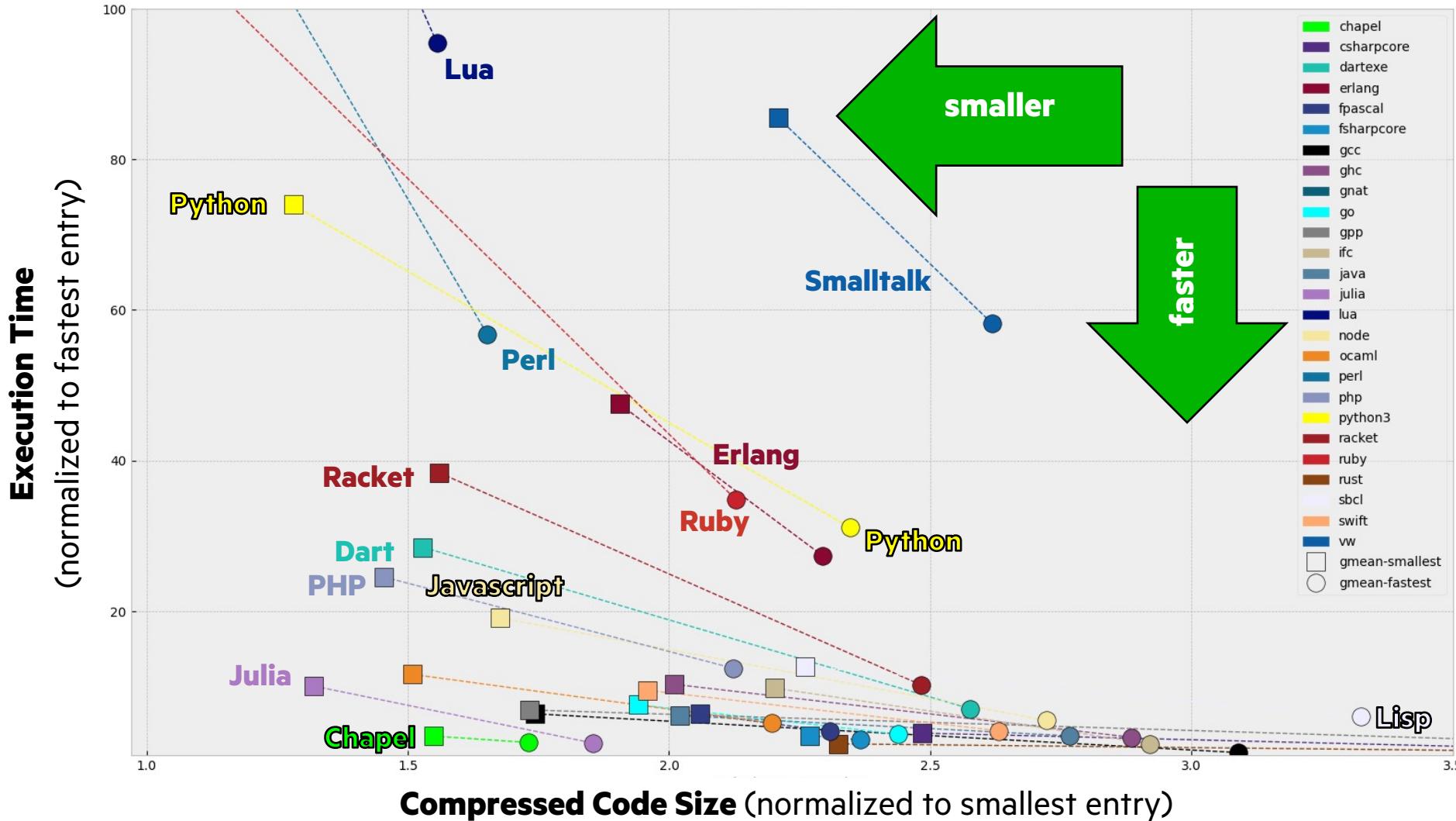
OUTLINE

- What is Chapel, and Why?
- Chapel Benchmarks and Apps
- Intro to Chapel, by Example
- Applications of Chapel
- Wrap-up



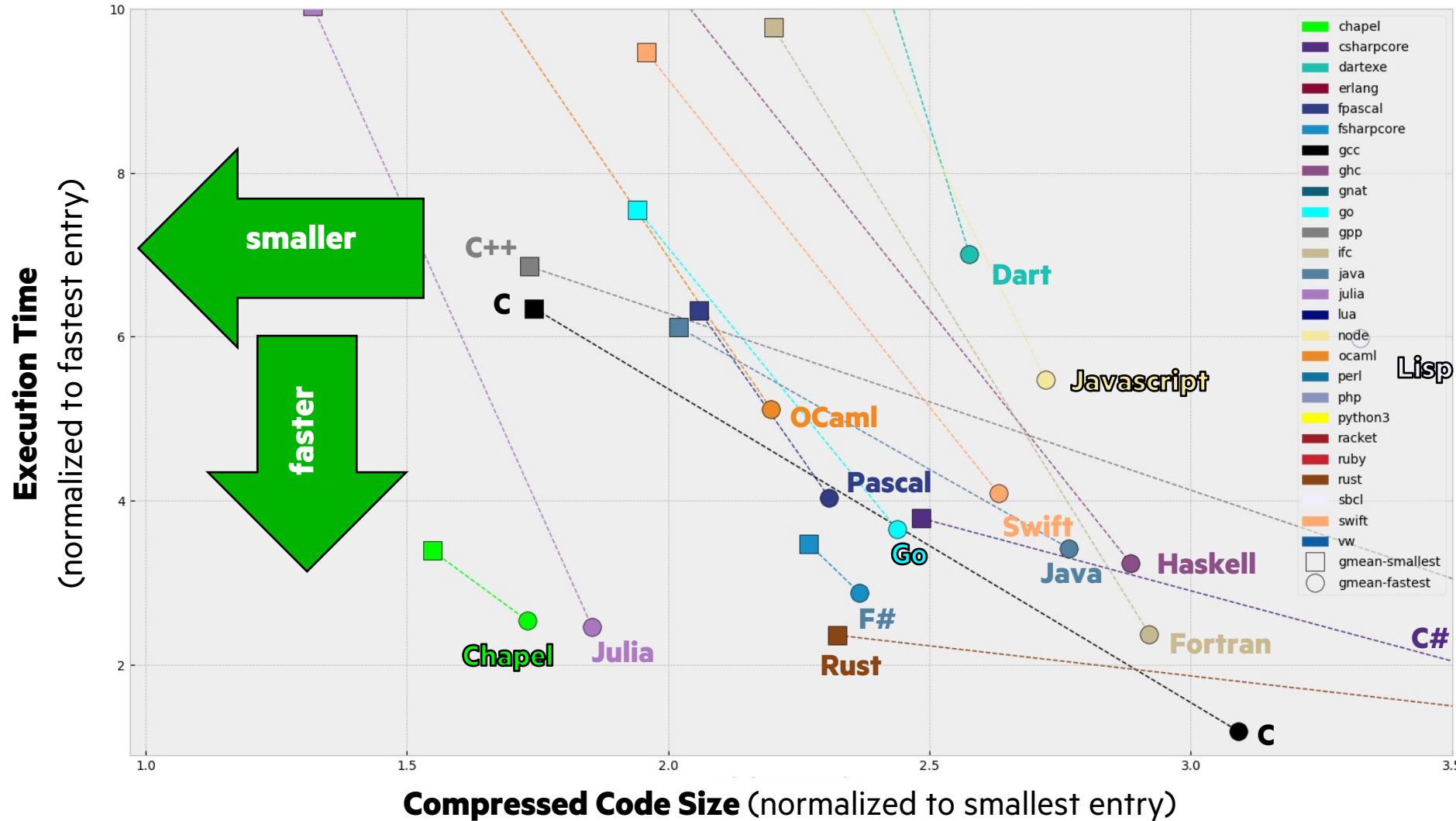
CHAPEL BENCHMARKS AND APPS

FOR DESKTOP BENCHMARKS, CHAPEL IS COMPACT AND FAST



[plot generated by summarizing data from <https://benchmarksgame-team.pages.debian.net/benchmarksgame/index.html> as of Feb 8, 2023]

FOR DESKTOP BENCHMARKS, CHAPEL IS COMPACT AND FAST (ZOOMED)



[plot generated by summarizing data from <https://benchmarksgame-team.pages.debian.net/benchmarksgame/index.html> as of Feb 8, 2023]

HPC BENCHMARKS: CONVENTIONAL APPROACHES VS. CHAPEL

STREAM TRIAD: C + MPI + OPENMP

```

/*include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;
    MPI_Comm_size(comm, &commSize);
    MPI_Comm_rank(comm, &myRank);

    rv = HPCC_Stream(params, 0 == myRank);
    MPI_Reduce(&rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm);

    return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
    register int j;
    double scalar;
    VectorSize = HPCC_LocalVectorSize( params, 3, sizeof(double), 0 );
    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );

    if (doIO) {
        #ifdef _OPENMP
        #pragma omp parallel
        #endif
        for (j=0; j < VectorSize; j++)
            a[j] = b[j];
        scalar = 3.0;
        #ifdef _OPENMP
        #pragma omp parallel
        #endif
        for (j=0; j < VectorSize; j++)
            c[j] = a[j] + scalar * b[j];
    }
    else {
        for (j=0; j < VectorSize; j++)
            c[j] = a[j] + alpha * b[j];
    }

    HPCC_free(c);
    HPCC_free(b);
    HPCC_free(a);
    return 0;
}

```

use **BlockDist**;

```

config const n = 1_000_000,
      alpha = 0.01;

const Dom = Block.createDomain({1..n});

var A, B, C: [Dom] real;

B = 2.0;
C = 1.0;

A = B + alpha * C;

```

HPCC RA: MPI KERNEL

```

/* Perform updates to main table. The scalar equivalent is:
   for (i=0;i<RASize;i++) Ra[i] = Ra[i] - 27.0*Ra[i]*Ra[i] + 0.7*POLY[i];
   MPI_Irecv(iLocalRaBuffer, localBuffSize, tparams.dtyped4,
             MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, iinreq);
   while (i < Sentsize);
   /* receive message */
   MPI_DemandRecv(shave_done, &status);
   if (&status.MPI_TAG == UPDATE_TAG) {
       if (shave_index >= 0) {
           if (shave_index < localBuffSize) {
               if (shave_index < localBuffSize - 1) {
                   imsg = LocalRecvbuffer[localBuffSize - 1];
                   locloffset = (imsg - localBuffStart) * globalOffset;
                   if (locloffset < localBuffStart)
                       locloffset = (imsg - localBuffStart) * globalOffset;
                   MPI_TableLocOffSet = ~imsg;
               }
               if (shave_index >= localBuffSize) {
                   if (shave_index < localBuffSize + localBuffSize) {
                       if (shave_index < localBuffSize + localBuffSize - 1) {
                           imsg = LocalRecvbuffer[localBuffSize];
                           locloffset = (imsg - localBuffStart) * globalOffset;
                           if (locloffset < localBuffStart)
                               locloffset = (imsg - localBuffStart) * globalOffset;
                           MPI_TableLocOffSet = ~imsg;
                       }
                   }
               }
           }
       }
       if (&status.MPI_TAG == FINISHED_TAG) {
           if (shave_index < localBuffSize) {
               if (shave_index < localBuffSize - 1) {
                   imsg = LocalRecvbuffer[localBuffSize - 1];
                   locloffset = (imsg - localBuffStart) * globalOffset;
                   if (locloffset < localBuffStart)
                       locloffset = (imsg - localBuffStart) * globalOffset;
                   MPI_TableLocOffSet = ~imsg;
               }
           }
       }
   }
   while (shave_index < NumberReceived) {
       if (pendingUpdates < maxPendingUpdates) {
           Ra = Ra << 1;
           if (Ra < 0) Ra = (raIndex + 1) * Ra;
           if (Ra < 0) Ra = (raIndex + 1) * Ra;
           if (GlobalOffset < tparams.Remainder) /
               tparams.GlobalTableSize;
           if (GlobalOffset == tparams.Remainder) /
               tparams.GlobalTableSize;
           if (GlobalOffset > tparams.Remainder - 1) -
               tparams.GlobalTableSize;
           while (have_done && NumberReceived > 0);
           MPI_Waitall();
       }
       else {
           if (shave_index < localBuffSize) {
               if (shave_index < localBuffSize - 1) {
                   imsg = LocalRecvbuffer[localBuffSize - 1];
                   locloffset = (imsg - localBuffStart) * globalOffset;
                   if (locloffset < localBuffStart)
                       locloffset = (imsg - localBuffStart) * globalOffset;
                   MPI_Irecv(iLocalRaBuffer, localBuffSize, tparams.dtyped4,
                             MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, iinreq);
                   MPI_Waitall();
               }
           }
       }
   }
}

```

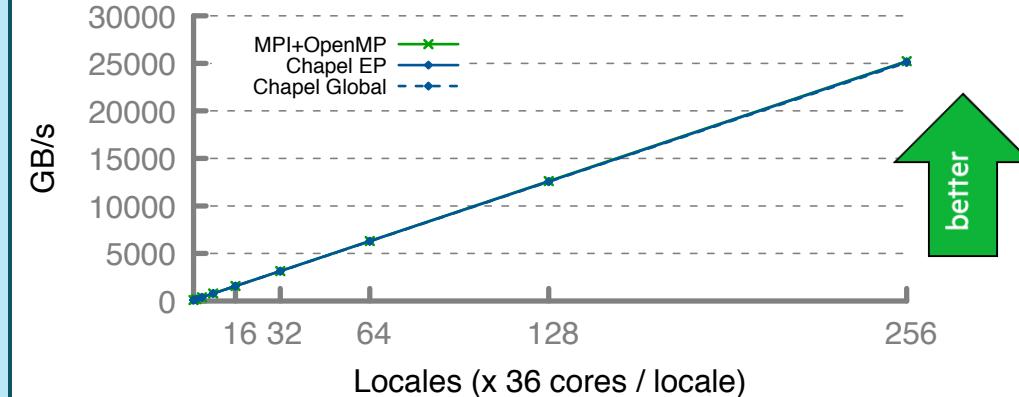
```

forall (_, r) in zip(Updates, RAStream()) do
    T[r & indexMask].xor(r);

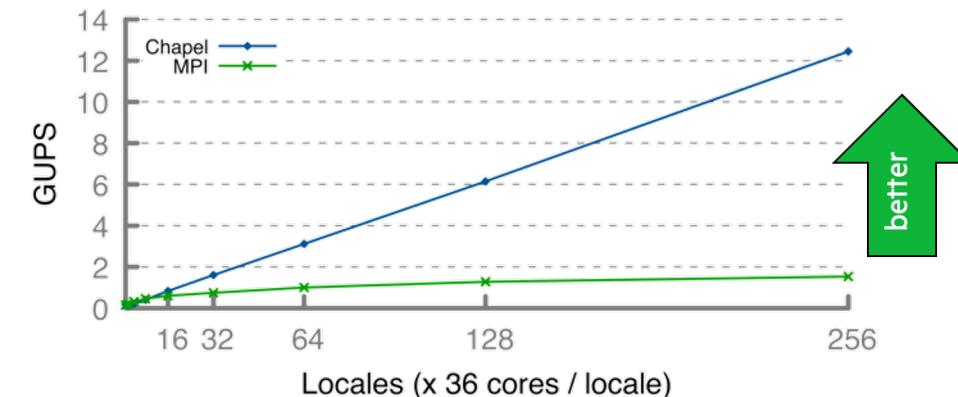
```

72

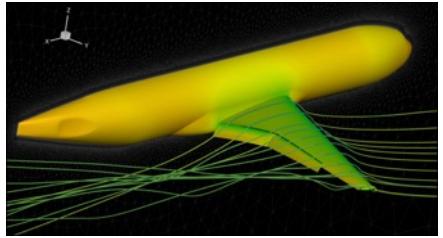
STREAM Performance (GB/s)



RA Performance (GUPS)

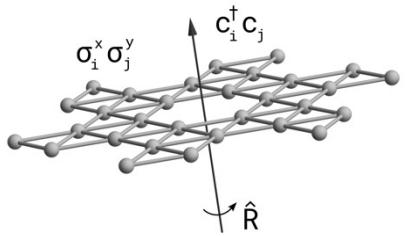


APPLICATIONS OF CHAPEL



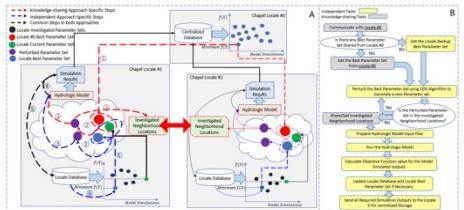
CHAMPS: 3D Unstructured CFD

Laurendeau, Bourgault-Côté, Parenteau, Plante, et al.
École Polytechnique Montréal



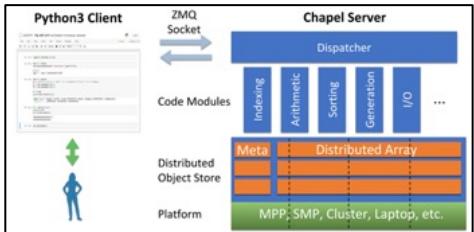
Lattice-Symmetries: a Quantum Many-Body Toolbox

Tom Westerhout
Radboud University



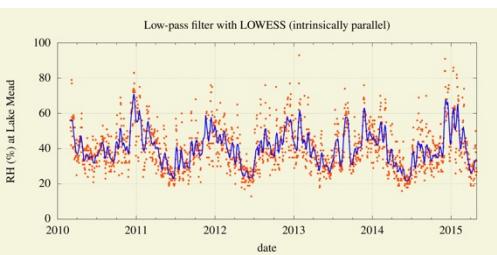
Chapel-based Hydrological Model Calibration

Marjan Asgari et al.
University of Guelph



Arkouda: Interactive Data Science at Massive Scale

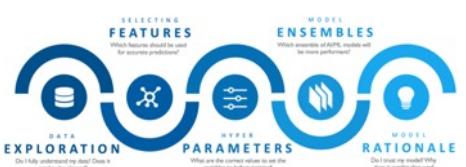
Mike Merrill, Bill Reus, et al.
U.S. DoD



Desk dot chpl: Utilities for Environmental Eng.

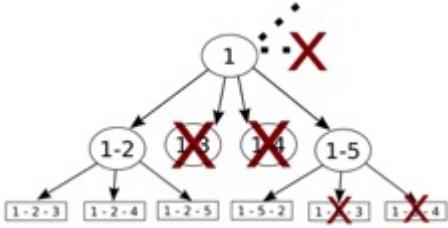
Nelson Luis Dias

The Federal University of Paraná, Brazil



CrayAI HyperParameter Optimization (HPO)

Ben Albrecht et al.
Cray Inc. / HPE



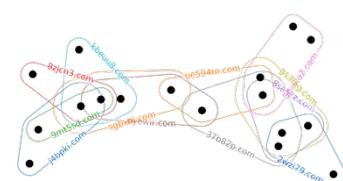
ChOp: Chapel-based Optimization

T. Carneiro, G. Helbecque, N. Melab, et al.
INRIA, IMEC, et al.



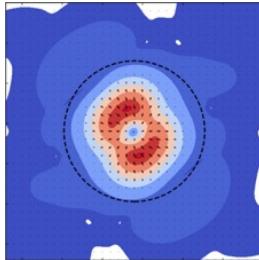
RapidQ: Mapping Coral Biodiversity

Rebecca Green, Helen Fox, Scott Bachman, et al.
The Coral Reef Alliance



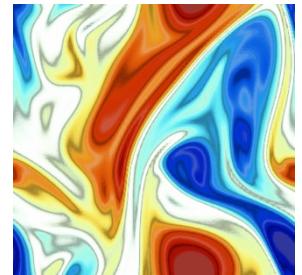
CHGL: Chapel Hypergraph Library

Louis Jenkins, Cliff Joslyn, Jesun Firoz, et al.
PNNL



ChplUltra: Simulating Ultralight Dark Matter

Nikhil Padmanabhan, J. Luna Zagorac, et al.
Yale University et al.



ChapQG: Layered Quasigeostrophic CFD

Ian Grooms and Scott Bachman
University of Colorado, Boulder et al.



Your Application Here?

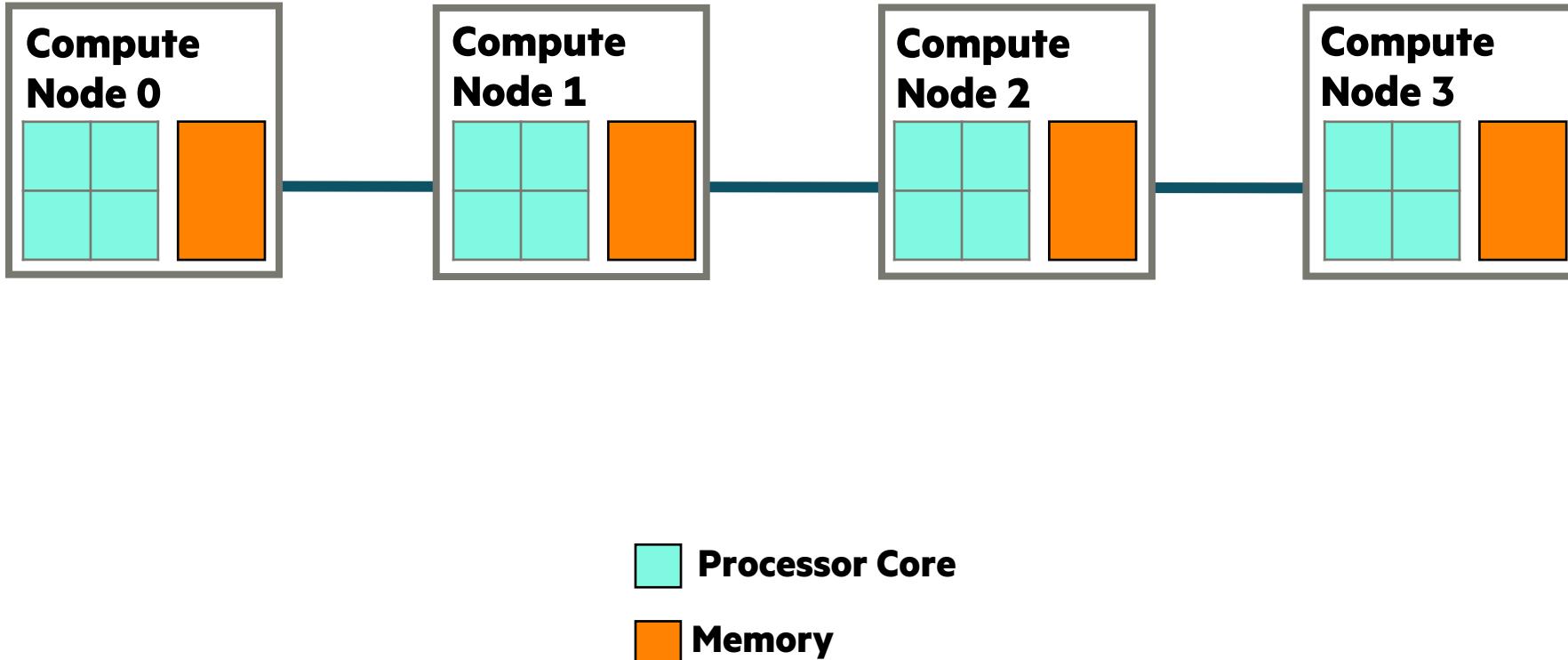
(Images provided by their respective teams and used with permission)

The background consists of several large, semi-transparent white circles of varying sizes that overlap each other, creating a sense of depth and motion.

INTRODUCTION TO CHAPEL, BY EXAMPLE

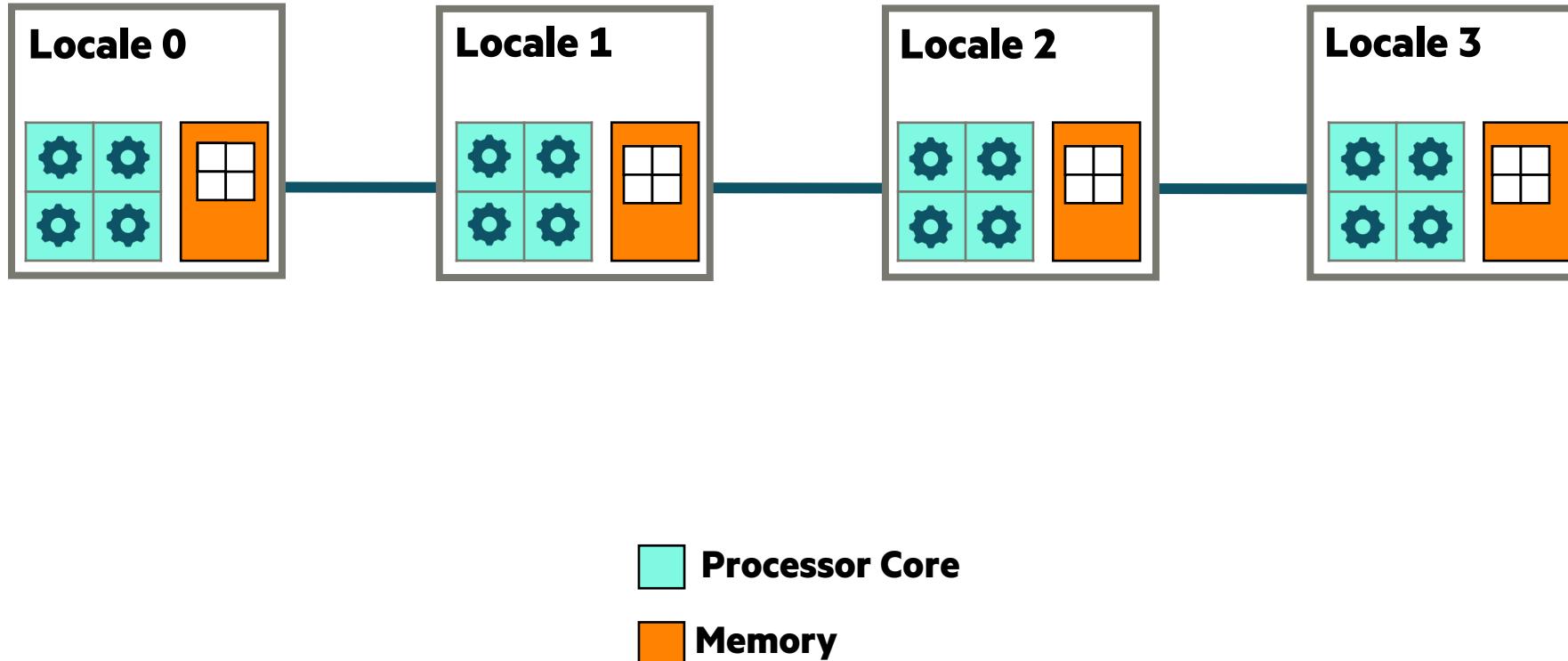
LOCALES IN CHAPEL

- In Chapel, a *locale* refers to a compute resource with...
 - processors, so it can run tasks
 - memory, so it can store variables
- For now, think of each compute node as being a locale



KEY CONCERNS FOR SCALABLE PARALLEL COMPUTING

- 1. parallelism:** What tasks should run simultaneously?
- 2. locality:** Where should tasks run? Where should data be allocated?



BASIC FEATURES FOR LOCALITY

basics-on.chpl

```
writeln("Hello from locale ", here.id);
```

```
var A: [1..2, 1..2] real;
```

```
on Locales[1] {
```

```
    var B: [1..2, 1..2] real;
```

```
    B = 2 * A;
```

```
}
```

All Chapel programs begin running as a single task on locale 0

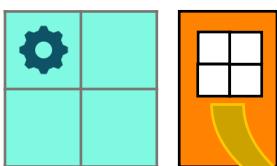
Variables are stored using the memory local to the current task

on-clauses move tasks to other locales

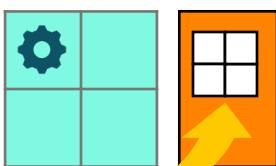
remote variables can be accessed directly

This is a serial, but distributed computation

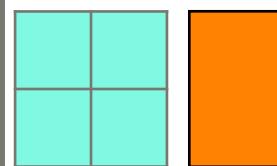
Locale 0



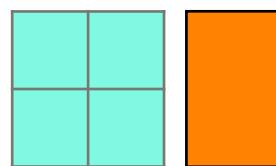
Locale 1



Locale 2



Locale 3



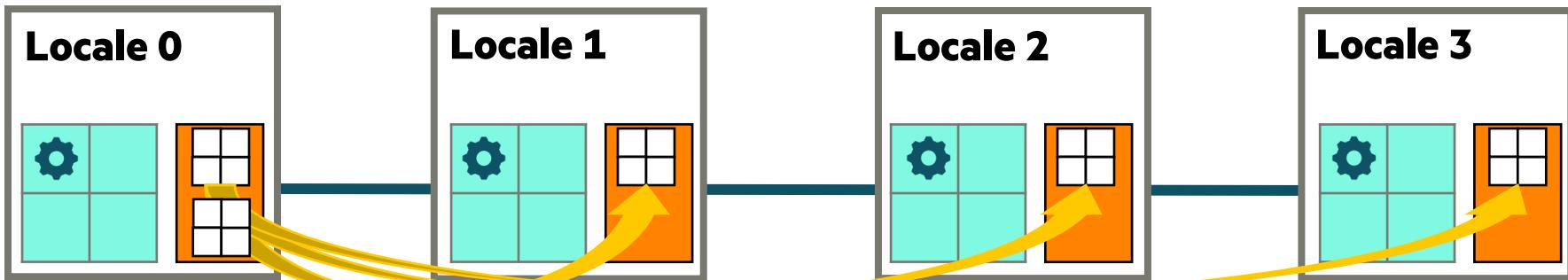
BASIC FEATURES FOR LOCALITY

basics-for.chpl

```
writeln("Hello from locale ", here.id);  
  
var A: [1..2, 1..2] real;  
  
for loc in Locales { ←  
  on loc {  
    var B = A;  
  }  
}
```

This loop will serially iterate over
the program's locales

This is also a serial, but distributed computation



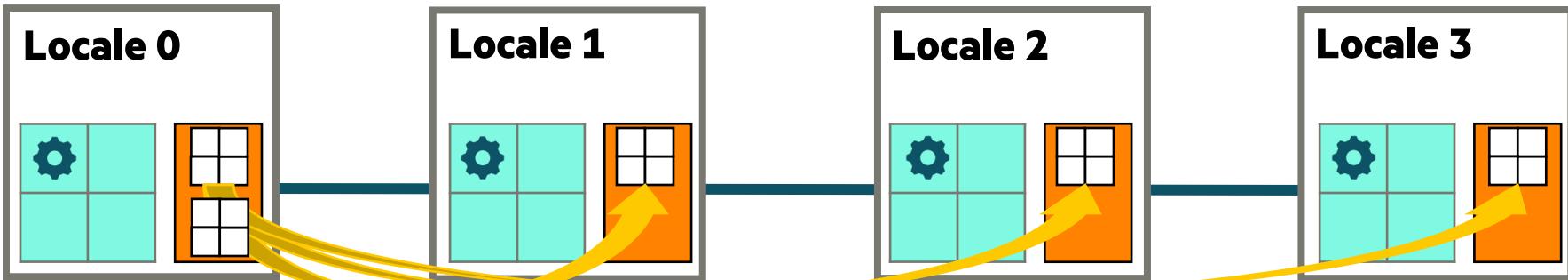
MIXING LOCALITY WITH TASK PARALLELISM

basics-coforall.chpl

```
writeln("Hello from locale ", here.id);  
  
var A: [1..2, 1..2] real;  
  
coforall loc in Locales { ←  
    on loc {  
        var B = A;  
    }  
}
```

The forall loop creates
a parallel task per iteration

This results in a parallel distributed computation



ARRAY-BASED PARALLELISM AND LOCALITY

basics-distarr.chpl

```
writeln("Hello from locale ", here.id);
```

```
var A: [1..2, 1..2] real;
```

```
use BlockDist;
```

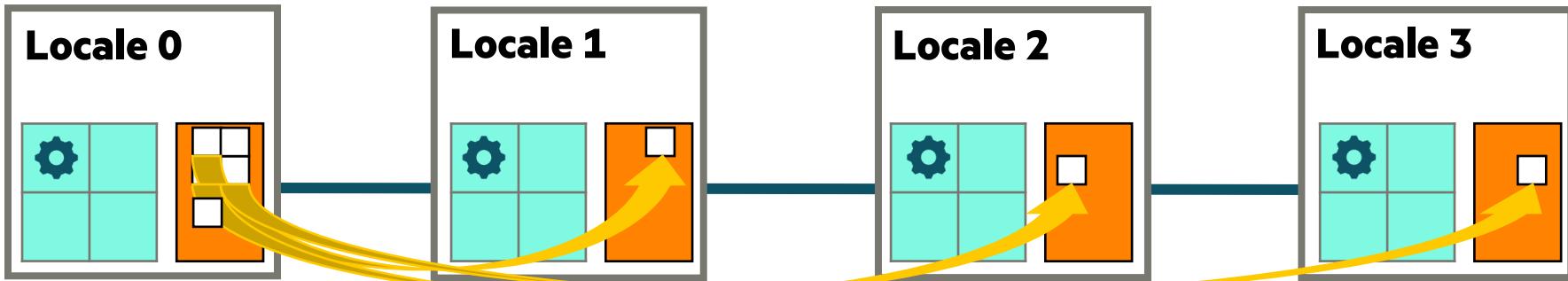
```
var D = Block.createDomain({1..2, 1..2});
```

```
var B: [D] real;
```

```
B = A;
```

Chapel also supports distributed domains (index sets) and arrays

They also result in parallel distributed computation

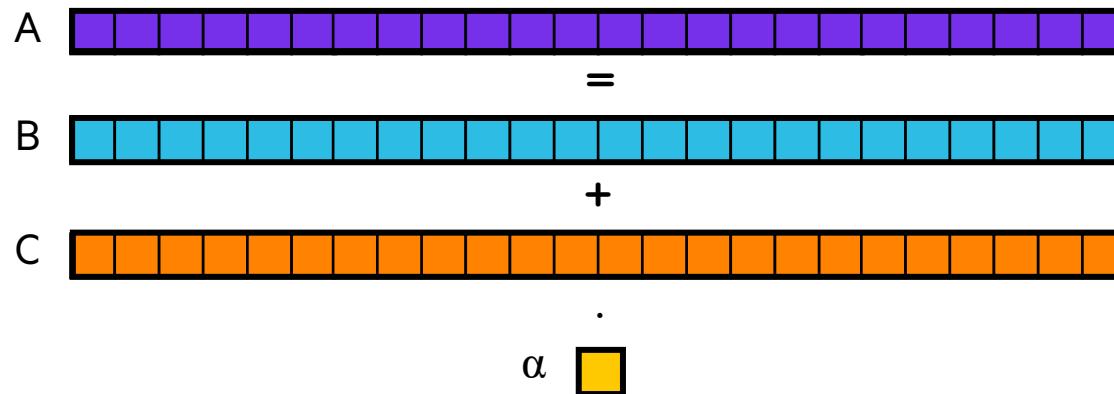


STREAM TRIAD: A TRIVIAL CASE OF PARALLELISM + LOCALITY

Given: n -element vectors A, B, C

Compute: $\forall i \in 1..n, A_i = B_i + \alpha \cdot C_i$

In pictures:

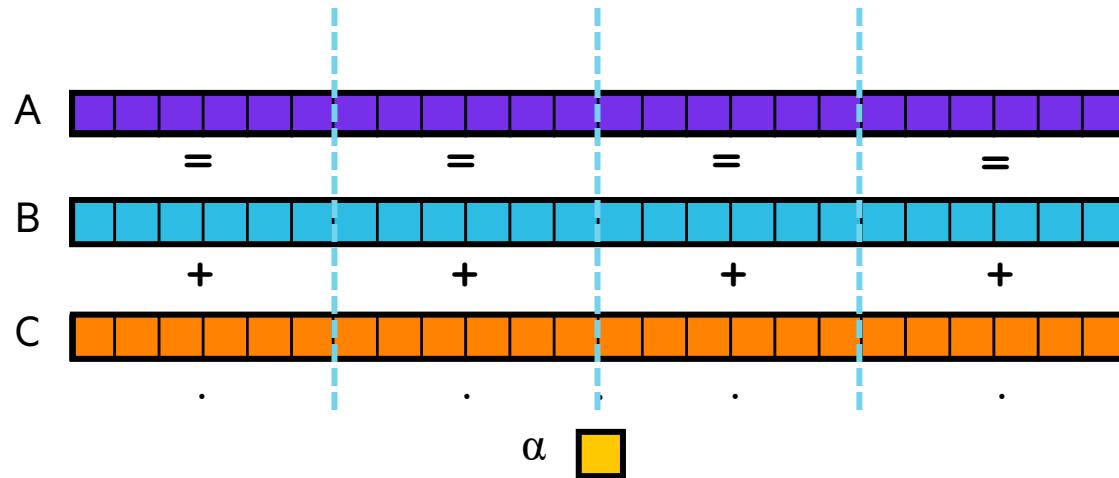


STREAM TRIAD: A TRIVIAL CASE OF PARALLELISM + LOCALITY

Given: n -element vectors A, B, C

Compute: $\forall i \in 1..n, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (shared memory / multicore):

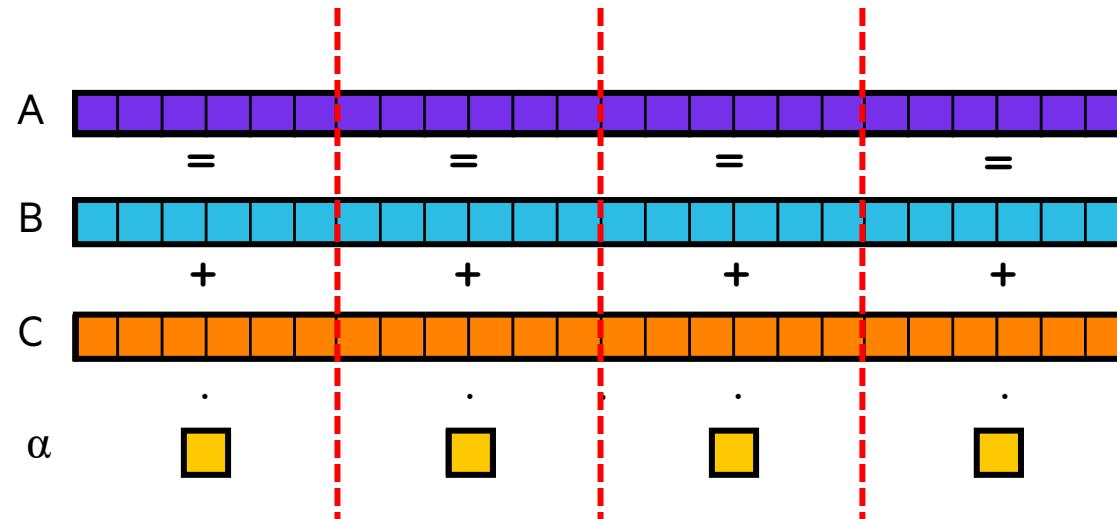


STREAM TRIAD: A TRIVIAL CASE OF PARALLELISM + LOCALITY

Given: n -element vectors A, B, C

Compute: $\forall i \in 1..n, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory):

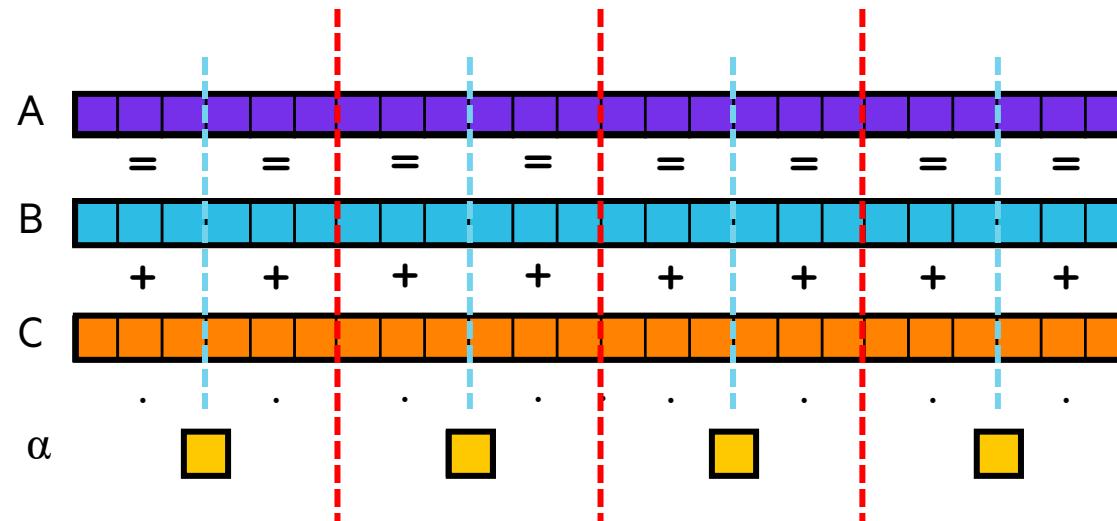


STREAM TRIAD: A TRIVIAL CASE OF PARALLELISM + LOCALITY

Given: n -element vectors A, B, C

Compute: $\forall i \in 1..n, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore):



STREAM TRIAD: SHARED MEMORY

stream-ep.chpl

```
config const n = 1_000_000,  
      alpha = 0.01;
```

'config' declarations support
command-line overrides

```
$ chpl stream-ep.chpl  
$ ./stream-ep  
$ ./stream-ep --n=10 --alpha=3.0
```

compile the program

run with the default values

override those values

n 

α 



STREAM TRIAD: SHARED MEMORY

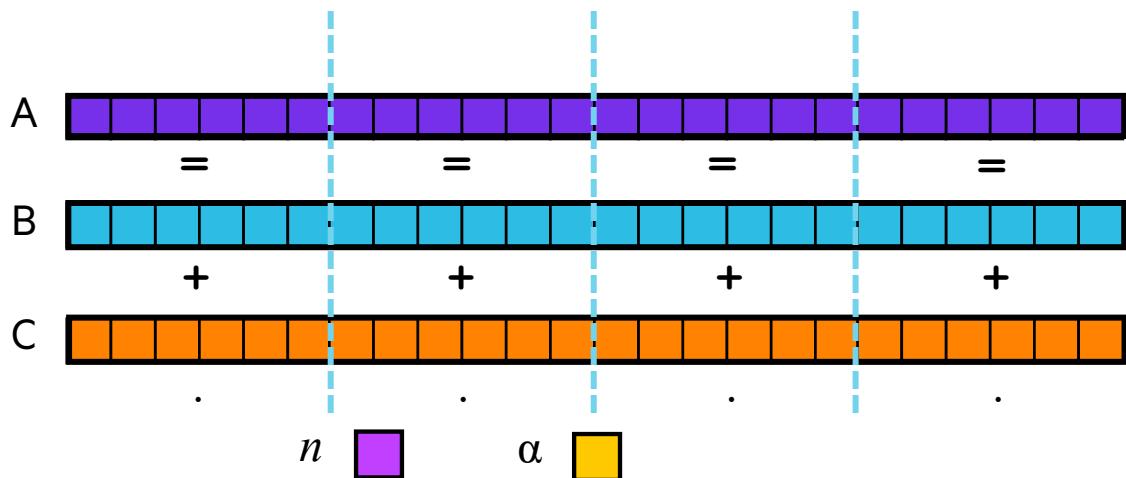
stream-ep.chpl

```
config const n = 1_000_000,  
      alpha = 0.01;
```

```
var A, B, C: [1..n] real;  
A = B + alpha * C;
```

declare three arrays of size ‘n’

whole-array operations result in
parallel computation



So far, this is simply a multi-core program

Nothing refers to remote locales,
explicitly or implicitly

STREAM TRIAD: DISTRIBUTED MEMORY (EP VERSION)

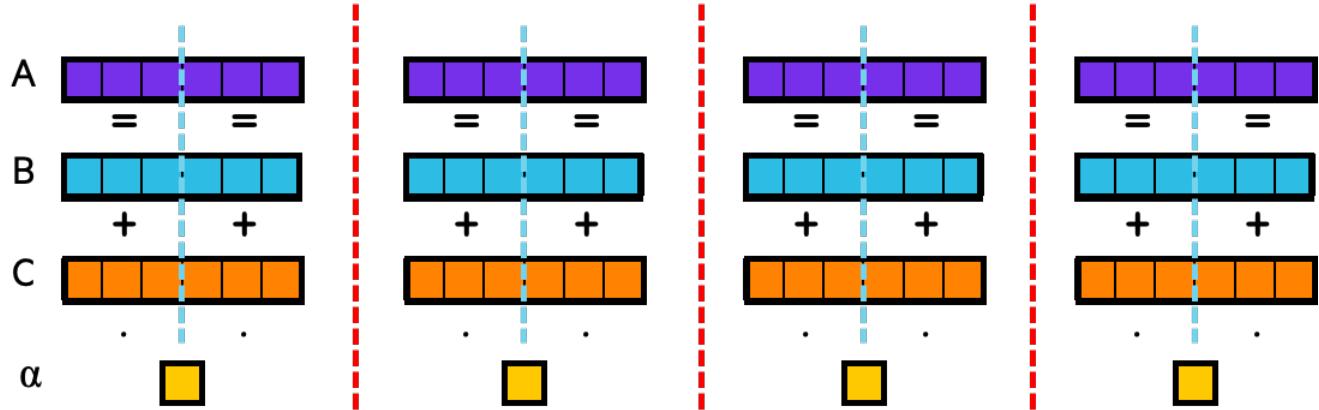
stream-ep.chpl

```
config const n = 1_000_000,  
        alpha = 0.01;  
  
coforall loc in Locales {  
    on loc {  
        var A, B, C: [1..n] real;  
        A = B + alpha * C;  
    }  
}
```

create a task per locale...

...running 'on' its locale

then run multi-core Stream
on local arrays, as before



STREAM TRIAD: DISTRIBUTED MEMORY (GLOBAL VERSION)

stream-glbl.chpl

```
config const n = 1_000_000,  
       alpha = 0.01;
```

```
use BlockDist;
```

```
const Dom = Block.createDomain({1..n});  
var A, B, C: [Dom] real;
```

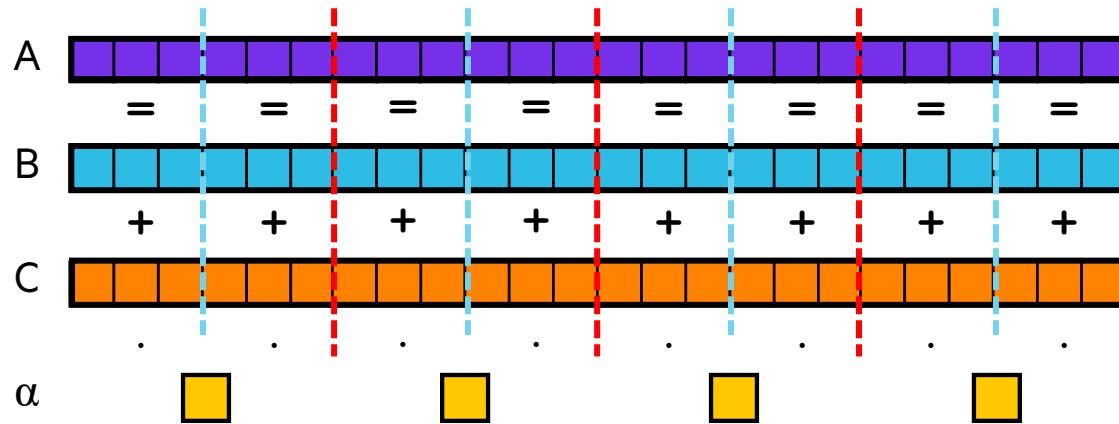
```
A = B + alpha * C;
```

'use' the standard block-distribution module

create a distributed domain (index set)...

...and distributed arrays

these whole-array operations
will use all cores on all locales



HPC BENCHMARKS: CONVENTIONAL APPROACHES VS. CHAPEL

STREAM TRIAD: C + MPI + OPENMP

```

#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StartStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM, 0, comm );
}

return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
register int j;
double scalar;

VectorSize = HPCC_LocalVectorSize( params, 3, sizeof(double) );

a = HPCC_XMALLOC( double, VectorSize );
b = HPCC_XMALLOC( double, VectorSize );
c = HPCC_XMALLOC( double, VectorSize );

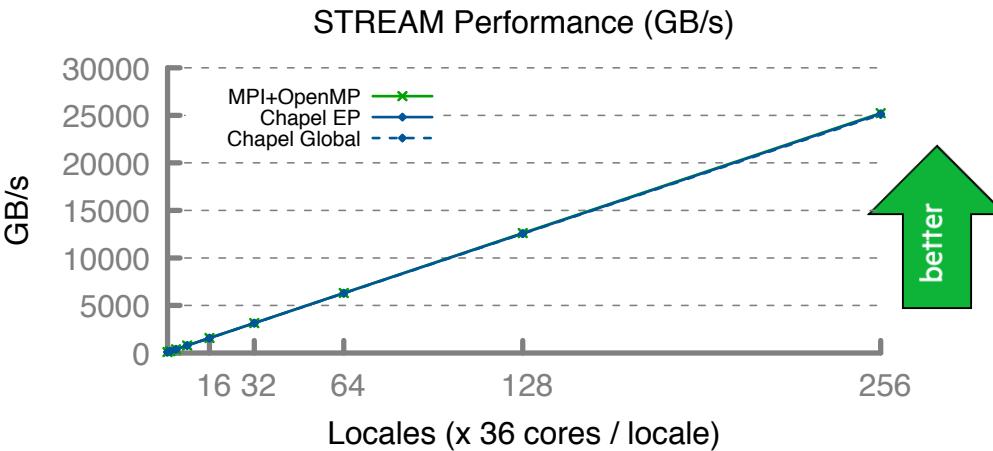
```

```
use BlockDist;

config const n = 1_000_000,
        alpha = 0.01;
const Dom = Block.createDomain({1..n});
var A, B, C: [Dom] real;

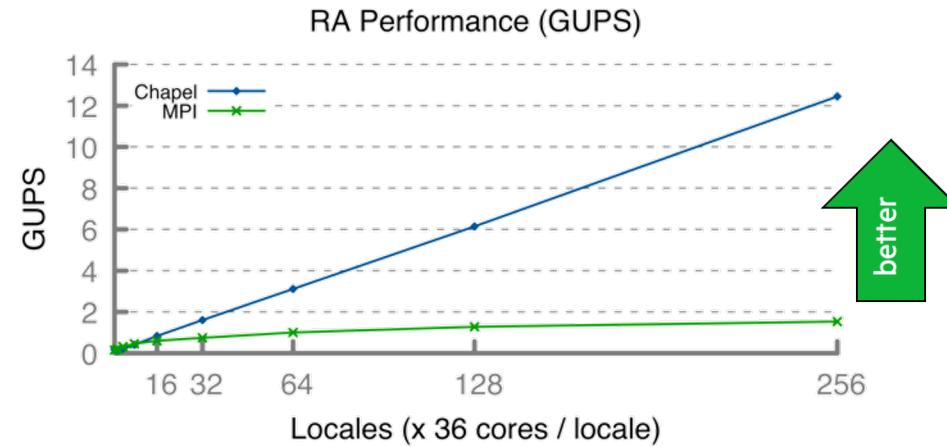
B = 2.0;
C = 1.0;

A = B + alpha * C;
```

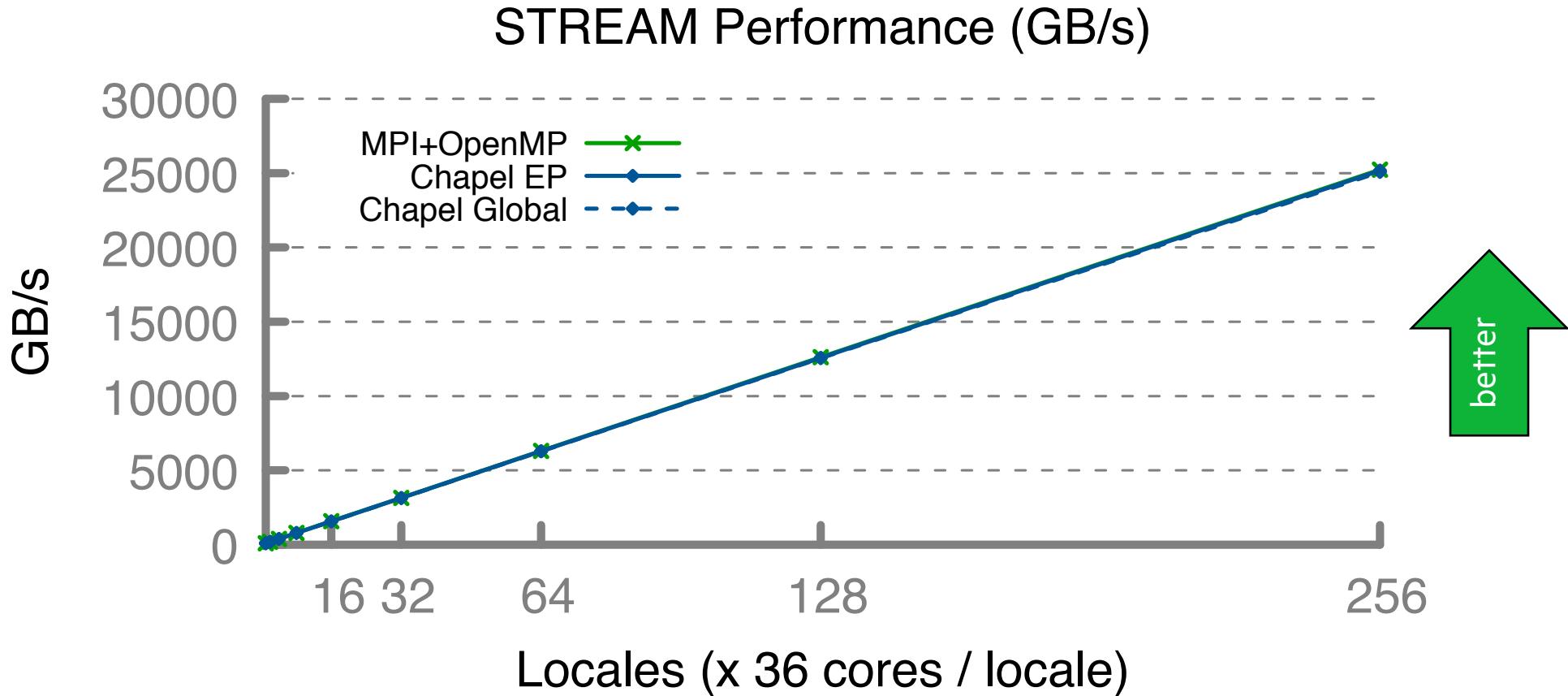


HPCC RA: MPI KERNEL

```
forall (_ , r) in zip(Updates, RASTream()) do
    T[r & indexMask].xor(r);
```

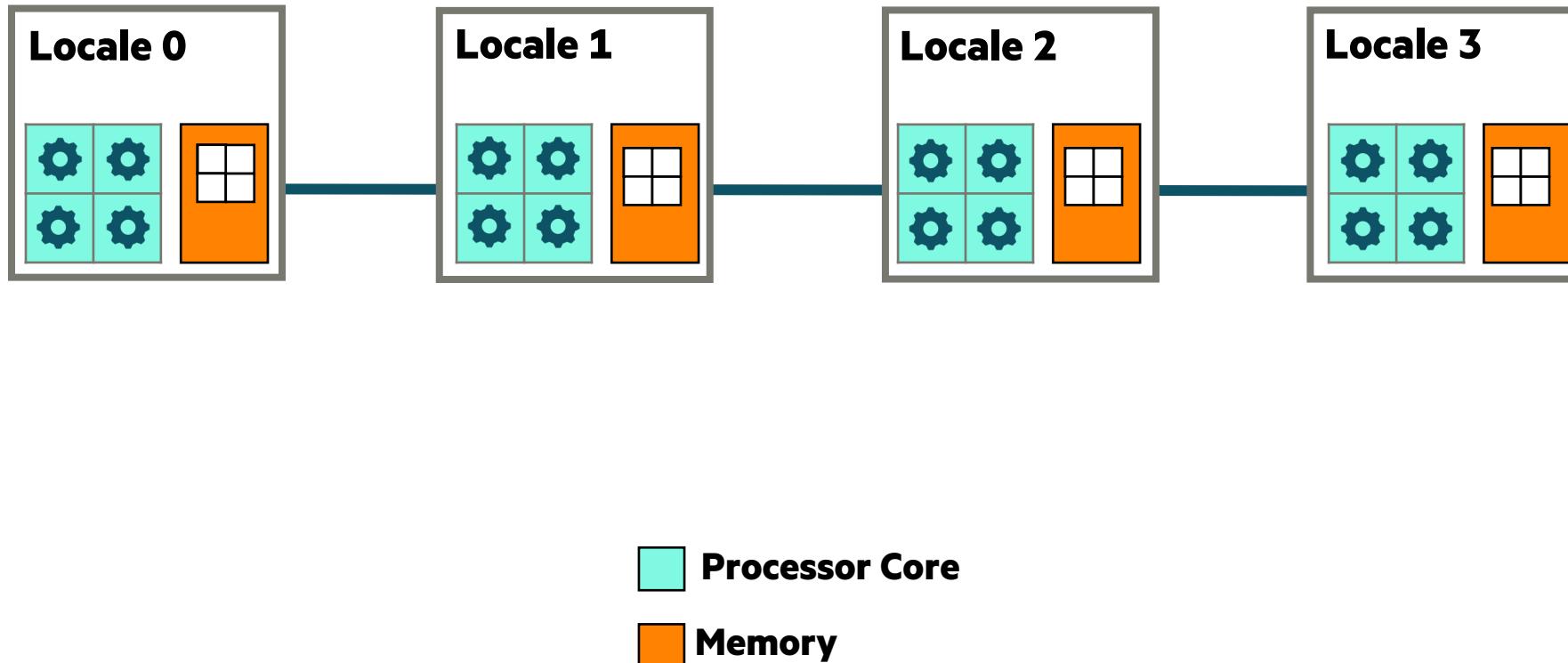


STREAM TRIAD: MPI + OPENMP VS. CHAPEL



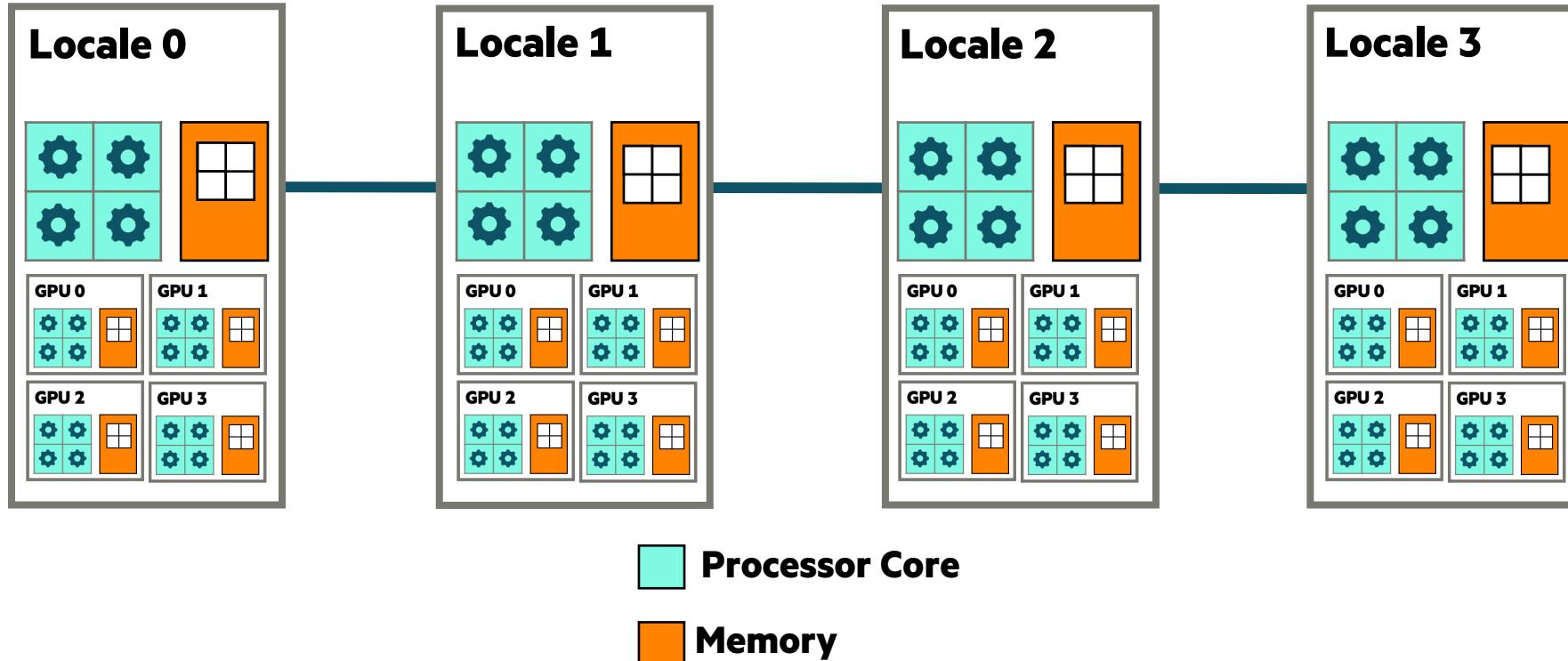
KEY CONCERN FOR SCALABLE PARALLEL COMPUTING

1. **parallelism:** What tasks should run simultaneously?
2. **locality:** Where should tasks run? Where should data be allocated?
 - complicating matters, compute nodes now often have GPUs with their own processors and memory



KEY CONCERN FOR SCALABLE PARALLEL COMPUTING

- parallelism:** What tasks should run simultaneously?
- locality:** Where should tasks run? Where should data be allocated?
 - complicating matters, compute nodes now often have GPUs with their own processors and memory
 - we represent these as *sub-locales* in Chapel



OAK RIDGE NATIONAL LABORATORY'S FRONTIER SUPERCOMPUTER



- 74 HPE Cray EX cabinets
- 9,408 AMD CPUs **37,632 AMD GPUs**
- 700 petabytes of storage capacity, peak write speeds of 5 terabytes per second using Cray ClusterStor storage system
- HPE Slingshot networking cables providing 100 GB/s network bandwidth.

TOP500

1

Built by HPE,
ORNL's Frontier
supercomputer
is #1 on the
TOP500.

1.1 exaflops of
performance.

TOP 500
The List.

GREEN500

2

Built by HPE,
ORNL's TDS and
full system are
ranked #2 & #6
on the Green500.

62.68 gigaflops/watt
power efficiency for
ORNL's TDS system,
52.23 gigaflops/watt
power efficiency for full
system.

The
GREEN
500

HPL-MxP

1

Built by HPE,
ORNL's Frontier
supercomputer
is #1 on the
HPL-MxP list.

7.9 exaflops on the
HPL-MxP benchmark
(formerly HPL-AI).

TOP 500
The List.

Source: May 30, 2022 [Top500](#) release, [HPL-MxP](#) mixed-precision benchmark (formerly HPL-AI).

STREAM TRIAD: DISTRIBUTED MEMORY, CPUS ONLY

stream-glbl.chpl

```
config const n = 1_000_000,  
        alpha = 0.01;  
  
use BlockDist;  
  
const Dom = Block.createDomain({1..n});  
var A, B, C: [Dom] real;  
  
A = B + alpha * C;
```

These programs are both CPU-only

Nothing refers to GPUs,
explicitly or implicitly

stream-ep.chpl

```
config const n = 1_000_000,  
        alpha = 0.01;  
  
coforall loc in Locales {  
    on loc {  
        var A, B, C: [1..n] real;  
        A = B + alpha * C;  
    }  
}
```

STREAM TRIAD: DISTRIBUTED MEMORY, GPUS ONLY

stream-ep.chpl

```
config const n = 1_000_000,  
        alpha = 0.01;  
  
coforall loc in Locales {  
    on loc {  
  
        coforall gpu in here.gpus do on gpu {  
            var A, B, C: [1..n] real;  
            A = B + alpha * C;  
        }  
    }  
}
```

Use a similar ‘coforall’ + ‘on’ idiom
to run a Triad concurrently
on each of this locale’s GPUs

This is a GPU-only program

Nothing other than coordination code
runs on the CPUs

STREAM TRIAD: DISTRIBUTED MEMORY, GPUS AND CPUS

stream-ep.chpl

```
config const n = 1_000_000,  
        alpha = 0.01;
```

```
coforall loc in Locales {
```

```
    on loc {
```

```
        cobegin {
```

```
            coforall gpu in here.gpus do on gpu {
```

```
                var A, B, C: [1..n] real;
```

```
                A = B + alpha * C;
```

```
}
```

```
{
```

```
    var A, B, C: [1..n] real;
```

```
    A = B + alpha * C;
```

```
}
```

```
}
```

```
}
```

'cobegin { ... }' creates a task per child statement

one task runs our multi-GPU triad

the other runs the multi-CPU triad

This program uses all CPUs and GPUs across all of our compute nodes

STREAM TRIAD: DISTRIBUTED MEMORY, GPUS AND CPUS (REFACTOR)

stream-ep.chpl

```
config const n = 1_000_000,
      alpha = 0.01;

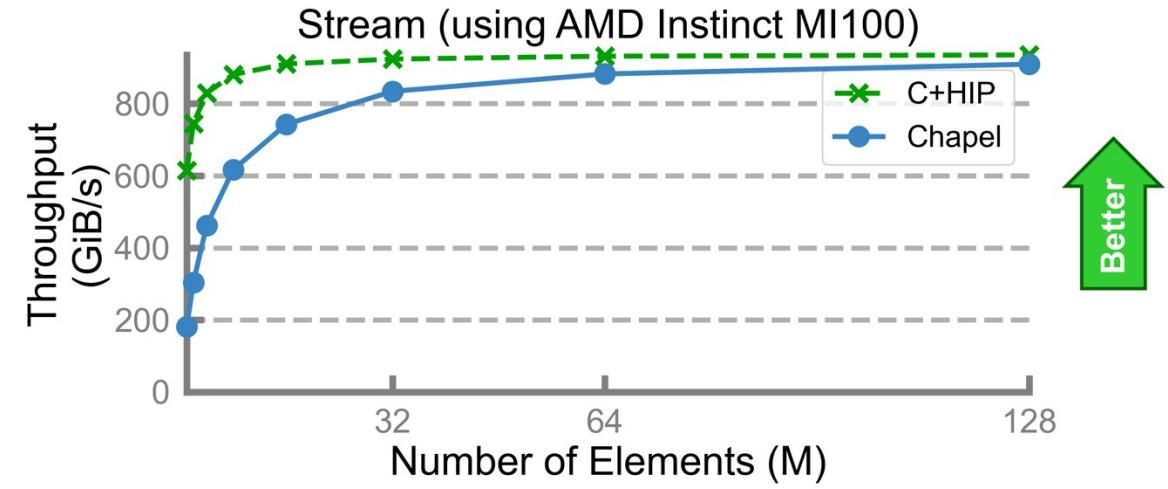
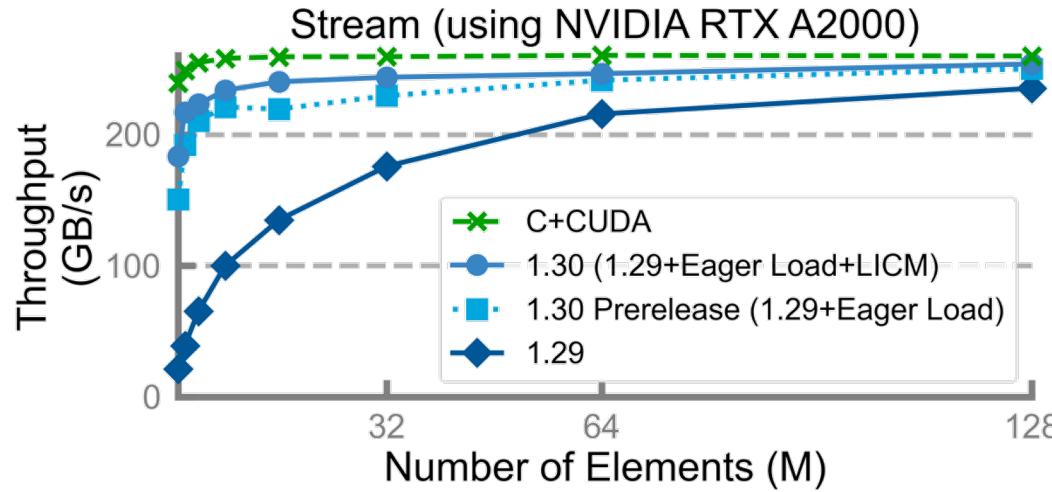
coforall loc in Locales {
    on loc {
        cobegin {
            coforall gpu in here.gpus do on gpu {
                runTriad();
            }
            runTriad();
        }
    }
}

proc runTriad() {
    var A, B, C: [1..n] real;
    A = B + alpha * C;
}
```

we can also refactor the repeated code into a procedure for re-use

the compiler creates CPU and GPU versions of this procedure

STREAM TRIAD: GPU PERFORMANCE VS. REFERENCE VERSIONS

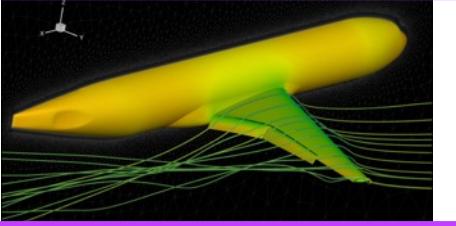


Performance vs. reference versions has become increasingly competitive over the past 4 months



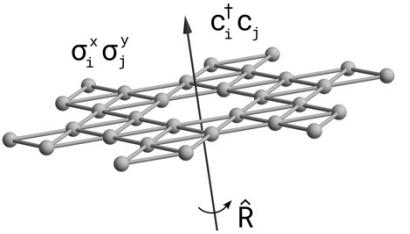
APPLICATIONS OF CHAPEL

APPLICATIONS OF CHAPEL



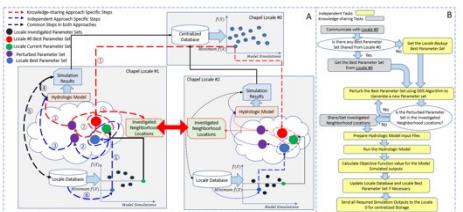
CHAMPS: 3D Unstructured CFD

Laurendeau, Bourgault-Côté, Parenteau, Plante, et al.
École Polytechnique Montréal

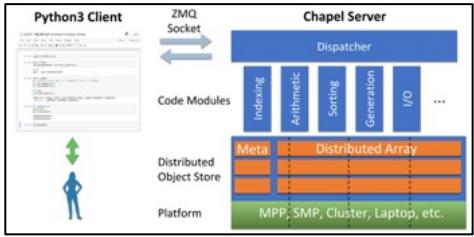


Lattice-Symmetries: a Quantum Many-Body Toolbox

Tom Westerhout
Radboud University

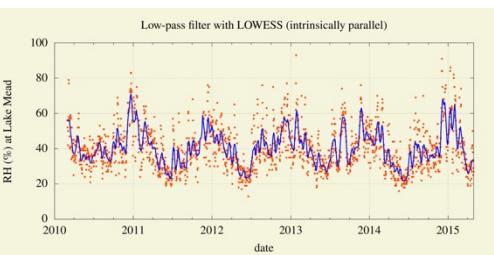


Chapel-based Hydrological Model Calibration
Marjan Asgari et al.
University of Guelph



Arkouda: Interactive Data Science at Massive Scale

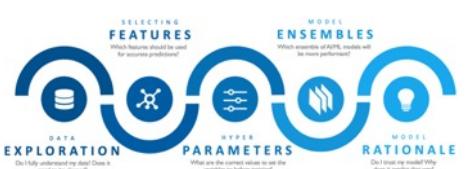
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U.S. DoD



Desk dot chpl: Utilities for Environmental Eng.

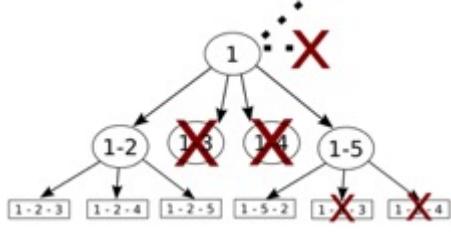
Nelson Luis Dias

The Federal University of Paraná, Brazil



CrayAI HyperParameter Optimization (HPO)

Ben Albrecht et al.
Cray Inc. / HPE



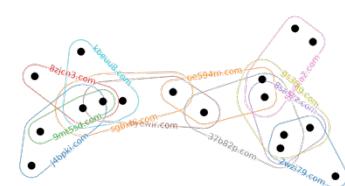
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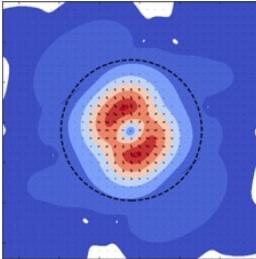
RapidQ: Mapping Coral Biodiversity

Rebecca Green, Helen Fox, Scott Bachman, et al.
The Coral Reef Alliance



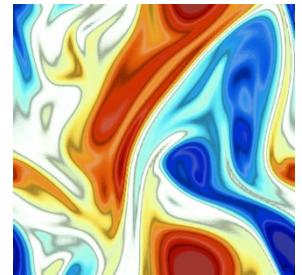
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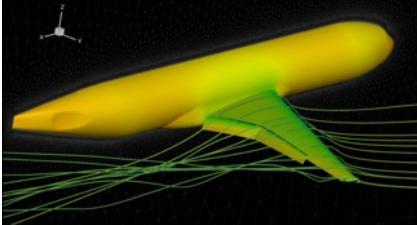
Your Application Here?

(Images provided by their respective teams and used with permission)

CHAMPS SUMMARY

What is it?

- 3D unstructured CFD framework for airplane simulation
- ~85k lines of Chapel written from scratch in ~3 years



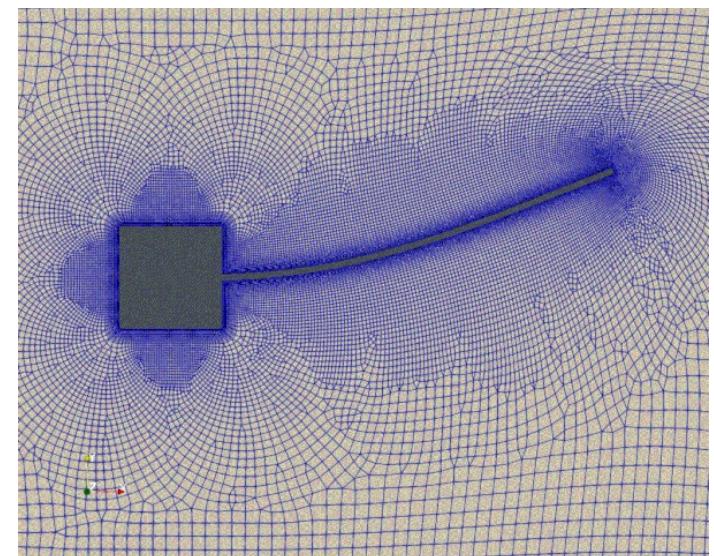
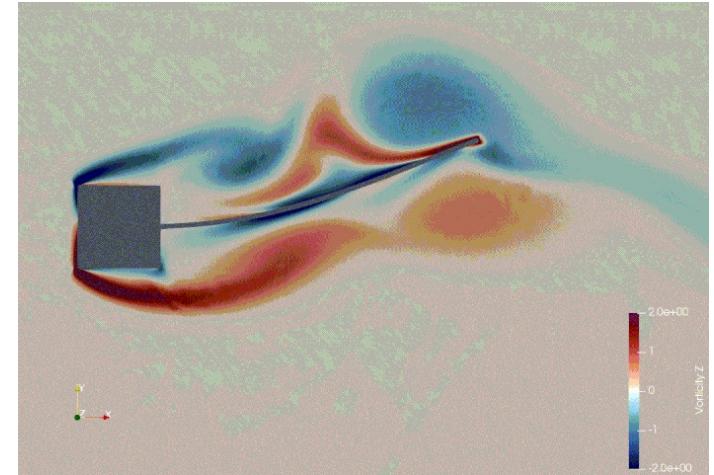
Who wrote it?

- Professor Éric Laurendeau's students + postdocs at Polytechnique Montreal



Why Chapel?

- performance and scalability competitive with MPI + C++
- students found it far more productive to use
- enabled them to compete with more established CFD centers



(images provided by the CHAMPS team and used with permission)

CHAMPS: EXCERPT FROM ÉRIC'S CHIUW 2021 KEYNOTE (TRANSCRIPT)

HPC Lessons From 30 Years of Practice in CFD Towards Aircraft Design and Analysis (June 4, 2021)

*"To show you what Chapel did in our lab... [our previous framework] ended up 120k lines. And my students said, 'We can't handle it anymore. It's too complex, we lost track of everything.' And today, they went **from 120k lines to 48k lines, so 3x less.***

*But the code is not 2D, it's 3D. And it's not structured, it's unstructured, which is way more complex. And it's multi-physics... **So, I've got industrial-type code in 48k lines.**"*

*"[Chapel] promotes the programming efficiency ... **We ask students at the master's degree to do stuff that would take 2 years and they do it in 3 months.** So, if you want to take a summer internship and you say, 'program a new turbulence model,' well they manage. And before, it was impossible to do."*

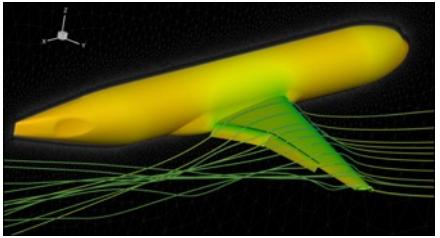
*"So, for me, this is like the proof of the benefit of Chapel, **plus the smiles I have on my students everyday in the lab because they love Chapel as well.** So that's the key, that's the takeaway."*

- Talk available online: https://youtu.be/wD-a_KyB8al?t=1904 (hyperlink jumps to the section quoted here)



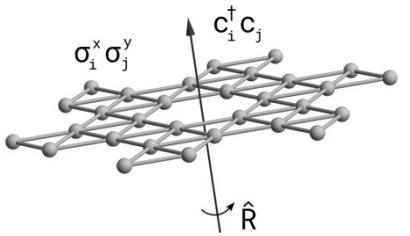
**POLYTECHNIQUE
MONTRÉAL**

APPLICATIONS OF CHAPEL



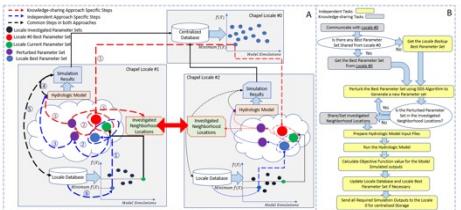
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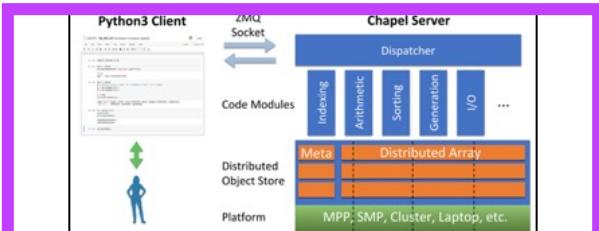
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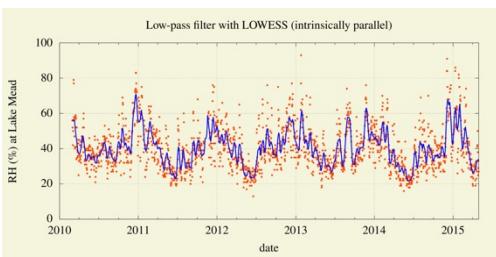
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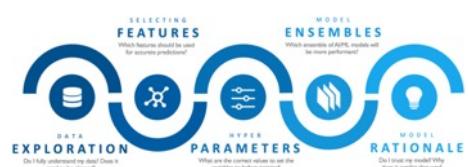
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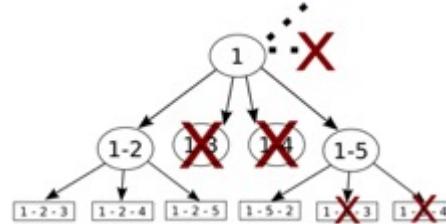
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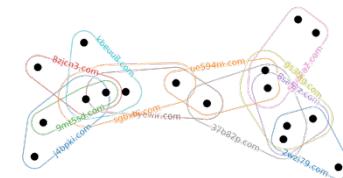
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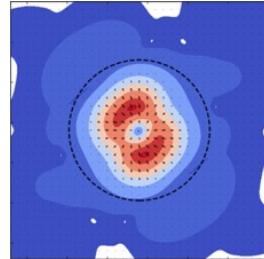
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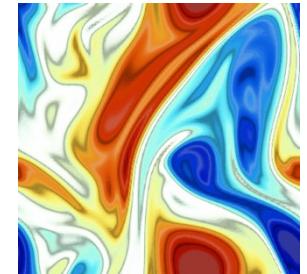
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University of Colorado, Boulder et al.



Your Application Here?

(Images provided by their respective teams and used with permission)

DATA SCIENCE IN PYTHON AT SCALE?

Motivation: Imagine you've got...

- ...HPC-scale data science problems to solve
- ...a bunch of Python programmers
- ...access to HPC systems

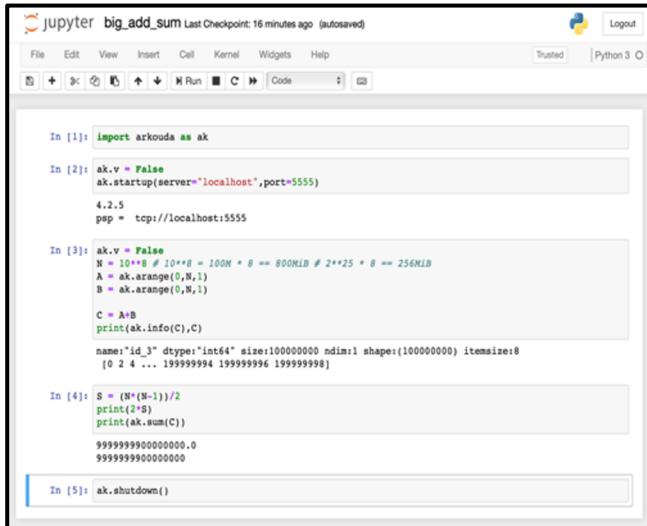


How will you leverage your Python programmers to get your work done?



ARKOUDA: A PYTHON FRAMEWORK FOR INTERACTIVE HPC

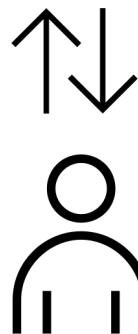
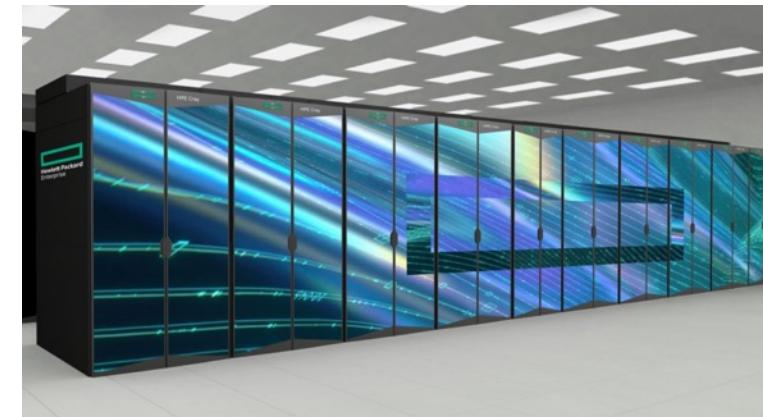
Arkouda Client (written in Python)



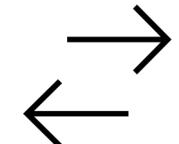
A screenshot of a Jupyter Notebook interface. The title bar says "jupyter big_add_sum Last Checkpoint: 16 minutes ago (autosaved)". The menu bar includes File, Edit, View, Insert, Cell, Kernel, Widgets, Help, Trusted, and Python 3. The code cell contains the following Python code:

```
In [1]: import arkouda as ak
In [2]: ak.v = False
ak.startup(server="localhost", port=5555)
4.2.5
psp = tcp://localhost:5555
In [3]: ak.v = True
N = 10**8 # 10**8 = 100M * 8 == 800MB # 2**25 * 8 == 256MB
A = ak.arange(0,N,1)
B = ak.arange(0,N,1)
C = A+B
print(ak.info(C),C)
name:id_3 dtype:int64 size:100000000 ndim:1 shape:(100000000) itemsize:8
[0 2 4 ... 199999994 199999996 199999998]
In [4]: S = (N*(N-1))/2
print(2*S)
print(ak.sum(C))
9999999900000000.0
9999999900000000.
In [5]: ak.shutdown()
```

Arkouda Server (written in Chapel)



User writes Python code in Jupyter,
making familiar NumPy/Pandas calls



ARKOUDA SUMMARY

What is it?

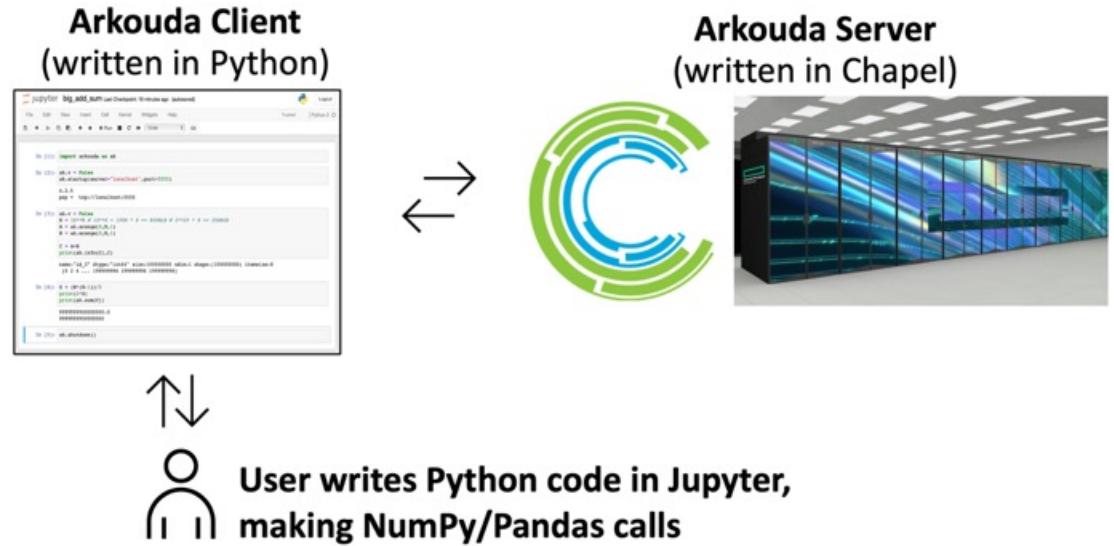
- A Python client-server framework supporting interactive supercomputing
 - Computes massive-scale results (TB-scale arrays) within the human thought loop (seconds to a few minutes)
 - Initial focus has been on a key subset of NumPy and Pandas for Data Science
- ~30k lines of Chapel + ~25k lines of Python, written since 2019
- Open-source: <https://github.com/Bears-R-Us/arkouda>

Who wrote it?

- Mike Merrill, Bill Reus, et al., US DoD

Why Chapel?

- close to Pythonic
 - enabled writing Arkouda rapidly
 - doesn't repel Python users who look under the hood
- achieved necessary performance and scalability
- ability to develop on laptop, deploy on supercomputer



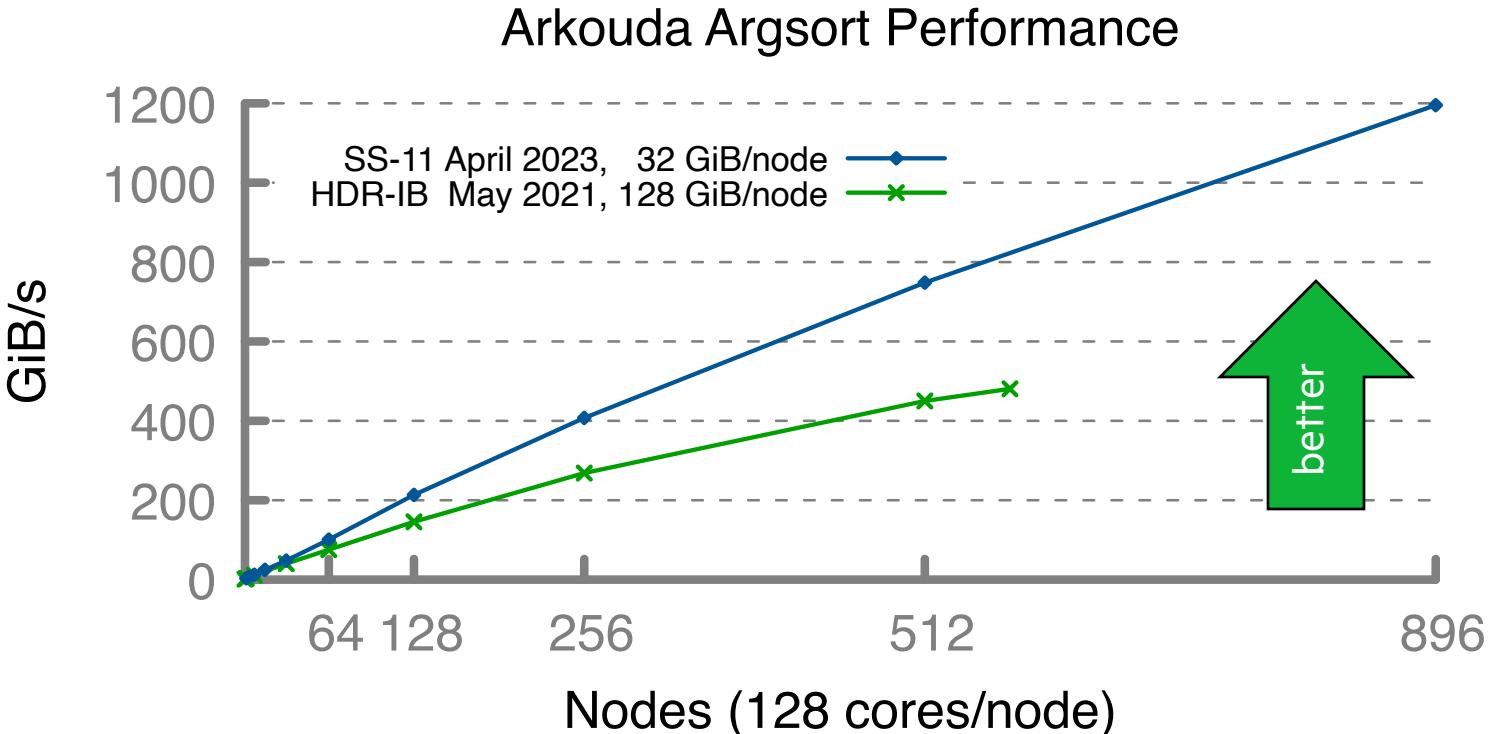
SCALABILITY OF ARKOUDA'S ARG SORT ROUTINE

HPE Cray EX (spring 2023)

- 114,688 cores of AMD Rome
- Slingshot-11 network (200 Gb/s)
- 28 TiB of 8-byte values
- 1200 GiB/s
 - 24 seconds elapsed time

HPE Apollo (summer 2021)

- 73,728 cores of AMD Rome
- HDR Infiniband network (100 Gb/s)
- 72 TiB of 8-byte values
- 480 GiB/s
 - 2.5 minutes elapsed time



A notable performance achievement in ~100 lines of Chapel



The background consists of numerous overlapping, slightly curved, light gray or white paper-like layers. These layers create a sense of depth and texture, resembling a stack of documents or a complex geometric pattern. The lighting is soft, highlighting the edges of the layers.

WRAP-UP

THE CHAPEL TEAM AT HPE



SUMMARY

Chapel is unique among programming languages

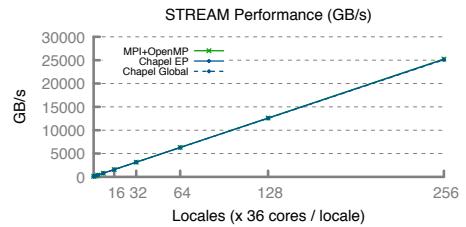
- built-in features for scalable parallel computing make it HPC-ready
- supports clean, concise code relative to conventional approaches
- ports and scales from laptops to supercomputers
- targets GPUs in a vendor-neutral manner

```
use BlockDist;

config const m = 1000,
        alpha = 3.0;
const Dom = {1..m} dimapped ...;
var A, B, C: [Dom] real;

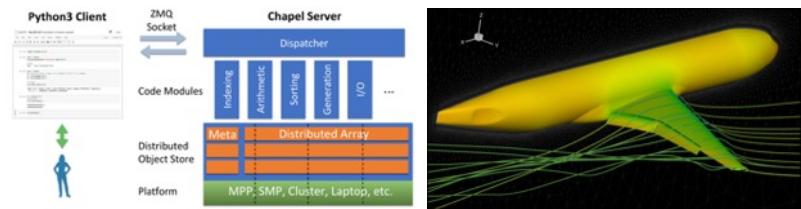
B = 2.0;
C = 1.0;

A = B + alpha * C;
```



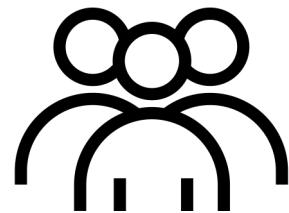
Chapel is being used for productive parallel computing at scale

- users are reaping its benefits in practical, cutting-edge applications
- applicable to domains as diverse as physical simulations and data science



If you or your users are interested in taking Chapel for a spin, let us know!

- we're happy to work with users and user groups to help ease the learning curve



COMING UP: CHIUW 2023



The Chapel Parallel Programming Language

CHIUW 2023

The 10th Annual
Chapel Implementers and Users Workshop

June 1–2, 2023
free and online in a virtual format

Home
What is Chapel?
What's New?
Blog
Upcoming Events
Job Opportunities

- **What?** The Chapel community's annual workshop
- **When?** June 1–2
 - one day of interactive programming
 - one day of presentations
- **Where?** Online
- **Cost?** Free

Details at: <https://chapel-lang.org/CHIUW2023.html>

CHAPEL RESOURCES

Chapel homepage: <https://chapel-lang.org>

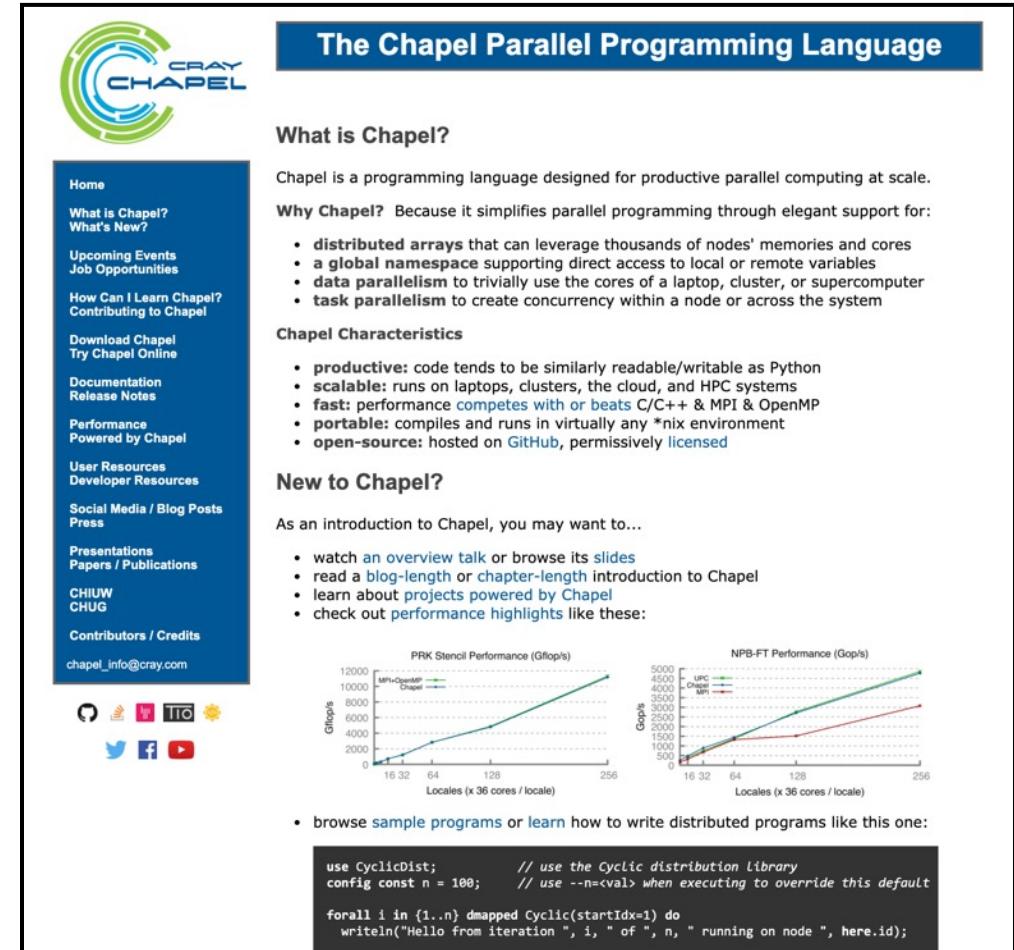
- (points to all other resources)

Social Media:

- Twitter: [@ChapelLanguage](#)
- Facebook: [@ChapelLanguage](#)
- YouTube: <http://www.youtube.com/c/ChapelParallelProgrammingLanguage>

Community Discussion / Support:

- Discourse: <https://chapel.discourse.group/>
- Gitter: <https://gitter.im/chapel-lang/chapel>
- Stack Overflow: <https://stackoverflow.com/questions/tagged/chapel>
- GitHub Issues: <https://github.com/chapel-lang/chapel/issues>



The screenshot shows the official Chapel website homepage. At the top right, a blue header bar reads "The Chapel Parallel Programming Language". To the left of the header is the Cray Chapel logo, which consists of a stylized green and blue circular icon followed by the word "CRAY" above "CHAPEL". The main content area has a dark blue sidebar on the left containing links such as Home, What is Chapel?, Upcoming Events, How Can I Learn Chapel?, Download Chapel, Documentation, User Resources, Social Media / Blog Posts, Presentations, CHI'16, CHUG, and Contributors / Credits. Below the sidebar is an email address: chapel_info@cray.com. To the right of the sidebar is a white main content area. At the top of this area is a section titled "What is Chapel?". Below it is a paragraph about Chapel's purpose: "Chapel is a programming language designed for productive parallel computing at scale." A section titled "Why Chapel?" follows, stating "Because it simplifies parallel programming through elegant support for:" and listing four bullet points: distributed arrays, a global namespace, data parallelism, and task parallelism. Further down are sections for "Chapel Characteristics" (with bullet points for productivity, scalability, performance, portability, and open-source), "New to Chapel?", and "As an introduction to Chapel, you may want to...". This section lists four items: watching an overview talk or slides, reading a blog-length or chapter-length introduction, learning about projects powered by Chapel, and checking out performance highlights. Two line graphs are shown: "PRK Stencil Performance (Gflop/s)" and "NPB-FT Performance (Gop/s)", both comparing MPI, OpenMP, UPC, and Chapel across 16, 32, 64, 128, and 256 locales. At the bottom of the page is a code snippet showing Chapel code for a cyclic distribution and iteration printing.

The Chapel Parallel Programming Language

What is Chapel?

Chapel is a programming language designed for productive parallel computing at scale.

Why Chapel? Because it simplifies parallel programming through elegant support for:

- distributed arrays that can leverage thousands of nodes' memories and cores
- a global namespace supporting direct access to local or remote variables
- data parallelism to trivially use the cores of a laptop, cluster, or supercomputer
- task parallelism to create concurrency within a node or across the system

Chapel Characteristics

- productive: code tends to be similarly readable/writable as Python
- scalable: runs on laptops, clusters, the cloud, and HPC systems
- fast: performance **competes with** or **beats** C/C++ & MPI & OpenMP
- portable: compiles and runs in virtually any *nix environment
- open-source: hosted on [GitHub](#), permissively licensed

New to Chapel?

As an introduction to Chapel, you may want to...

- watch an [overview talk](#) or browse its [slides](#)
- read a [blog-length](#) or [chapter-length](#) introduction to Chapel
- learn about [projects powered by Chapel](#)
- check out [performance highlights](#) like these:

PRK Stencil Performance (Gflop/s)

NPB-FT Performance (Gop/s)

```
use CyclicDist;           // use the Cyclic distribution Library
config const n = 100;      // use --n=<val> when executing to override this default
forall i in {1..n} dmapped Cyclic(startIdx=1) do
    writeln("Hello from iteration ", i, " of ", n, " running on node ", here.id);
```

SUMMARY

Chapel is unique among programming languages

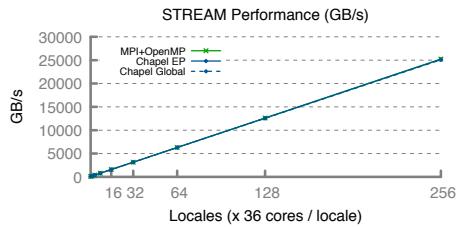
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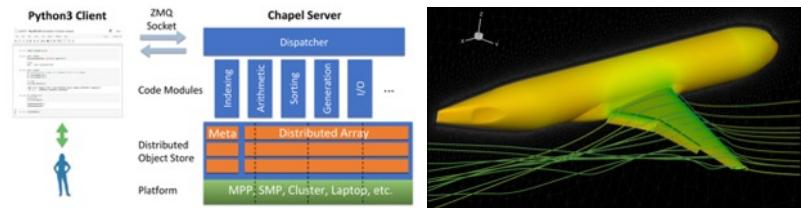
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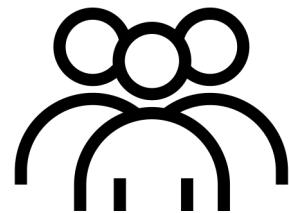
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THANK YOU

<https://chapel-lang.org>
@ChapelLanguage

